

Practice Problems from Section 3.6. Exercises 3.6.1–3.6.13 except for: 2, 10, 13.
 These *Practice Problems* are a sampling of the type of problems which could be on the exam.
 These *Practice Problem* are, in no way, meant as a comprehensive review for the exam.

Math is not a spectator sport.

Often we learn more from our failed attempts at a proof rather than immediately looking at the hints or reading a clean proof.

So give these problems a solid attempt before seeking help, e.g.:
 looking through: your notes, the textbook, the homework.

Since these problems are not to hand in, you may share:
 hints and/or an attempt at a solution for others to comment on.

Some hints are very generous. Do not except such generous hints on the exam.

1. **Theorem 1.** If $n \in \mathbb{N}$ with $n \geq 2$, then

$$\prod_{i=2}^n \frac{i^2 - 1}{i^2} = \frac{n + 1}{2n}.$$

1.1. Symbolically write Theorem 1. (Do not forget your quantifiers.)

1.2. Prove Theorem 1 using math induction..

o. Recall, the sum $\sum_{i=2}^5 a_i = a_2 + a_3 + a_4 + a_5$ while the product $\prod_{i=2}^5 a_i = a_2 \cdot a_3 \cdot a_4 \cdot a_5$

hint. The sum $\sum_{i=2}^{n+1} a_i = (a_2 + a_3 + \dots + a_n) + a_{n+1} = \left(\sum_{i=2}^n a_i\right) + a_{n+1}$. This idea is often used in the inductive step. Similarly, for the product $\prod_{i=2}^{n+1} a_i = (a_2 \cdot a_3 \cdot \dots \cdot a_n) \cdot a_{n+1} = \left(\prod_{i=2}^n a_i\right) \cdot a_{n+1}$.

sw. $(\forall n \in \mathbb{N}^{\geq 2}) \left[\prod_{i=2}^n \frac{i^2 - 1}{i^2} = \frac{n + 1}{2n} \right]$

Proof. We will show that for each natural number n with $n \geq 2$,

$$\prod_{i=2}^n \frac{i^2 - 1}{i^2} = \frac{n + 1}{2n} \tag{1.1}$$

by basic induction on the variable n .

For the base step, let $n = 2$. Since

$$\prod_{i=2}^2 \frac{i^2 - 1}{i^2} = \prod_{i=2}^2 \frac{i^2 - 1}{i^2} = \frac{2^2 - 1}{2^2} = \frac{3}{4} \quad \text{and} \quad \frac{n + 1}{2n} = \frac{2 + 1}{2(2)} = \frac{3}{4}.$$

the equation in (1.1) is true when $n = 2$. This completes the base step.

For the inductive step, fix $n \in \mathbb{N}^{\geq 2}$. We assume the inductive hypothesis, which is

$$\prod_{i=2}^n \frac{i^2 - 1}{i^2} = \frac{n + 1}{2n}. \tag{IH}$$

We will show the inductive conclusion, which is

$$\prod_{i=2}^{n+1} \frac{i^2 - 1}{i^2} = \frac{(n + 1) + 1}{2(n + 1)}. \tag{IC}$$

We know that (recall $\prod_{i=2}^{n+1} a_i = \left[\prod_{i=2}^n a_i \right] \cdot a_{n+1}$)

$$\prod_{i=2}^{n+1} \frac{i^2 - 1}{i^2} = \left[\prod_{i=2}^n \frac{i^2 - 1}{i^2} \right] \left[\frac{(n+1)^2 - 1}{(n+1)^2} \right].$$

and using the inductive hypothesis (IH) gives

$$= \left[\frac{n+1}{2n} \right] \left[\frac{(n+1)^2 - 1}{(n+1)^2} \right].$$

and now through algebra (note the algebraic cancellation in next step, which greatly simplifies the algebra to come - a good lesson to learn and practice)

$$\begin{aligned} &= \frac{(n+1)^2 - 1}{2n(n+1)} \\ &= \frac{n^2 + 2n}{2n(n+1)} \\ &= \frac{n}{n} \frac{(n+2)}{2(n+1)} \\ &= \frac{(n+1) + 1}{2(n+1)}. \end{aligned} \tag{1.2}$$

The calculations through (1.2) shows that the inductive conclusion (IC) holds. This completes the inductive step.

Thus the base step and inductive step hold. \square

2. Theorem 2. If $n \in \mathbb{N}$ with $n \geq 3$, then

$$2n + 1 < 2^n .$$

2.1. Symbolically write Theorem 2. (Do not forget your quantifiers.)

2.2. Prove Theorem 2 using math induction.

sw. $(\forall n \in \mathbb{N}^{\geq 3}) [2n + 1 < 2^n]$

Proof. We will show that if $n \in \mathbb{N}^{\geq 3}$, then

$$2n + 1 < 2^n \tag{2.1}$$

by basic induction on the variable n .

For the base step, let $n = 3$. Since

$$2n + 1 = 2(3) + 1 = 7 \quad \text{and} \quad 2^n = 2^3 = 8$$

and $7 < 8$, the equation in (2.1) is true when $n = 3$. This completes the base step.

For the inductive step, fix $n \in \mathbb{N}^{\geq 3}$. We assume the inductive hypothesis, which is

$$2n + 1 < 2^n. \tag{IH}$$

We will show the inductive conclusion, which is

$$2(n+1) + 1 < 2^{n+1}. \tag{IC}$$

From algebra we know

$$\begin{aligned} 2(n+1) + 1 &= 2n + 2 + 1 \\ &= (2n + 1) + 2 \end{aligned}$$

and by the (IH)

$$< (2^n) + 2$$

and since $3 \leq n$ we get $2 < 2^3 \leq 2^n$ so

$$\begin{aligned} &< 2^n + 2^n \\ &= 2^1 (2^n) \\ &= 2^{n+1} \end{aligned}$$

and so the inductive conclusion (IC) holds. This completes the inductive step.

Thus the base step and inductive step hold. \square

3. Theorem 3. If $n \in \mathbb{Z}$ with $n \geq 0$, then

$$3 \mid (n^3 + 2n).$$

3.1. Symbolically write Theorem 3. (Do not forget your quantifiers.)

3.2. Prove Theorem 3 using math induction.

hint. See book Proposition 4.4 (p178) and (starred) ER 3.4.8a (p181).

sw. $(\forall n \in \mathbb{Z}^{\geq 0}) [3 \mid (n^3 + 2n)]$

Proof. We will show that if $n \in \mathbb{Z}^{\geq 0}$, then

$$3 \mid (n^3 + 2n) \tag{3.1}$$

by basic induction on the variable n .

For the base step, let $n = 0$. Then

$$n^3 + 2n = 0^3 + 2(0) = 0.$$

Also

$$3 \mid 0$$

since $0 = 3j$ where $j = 0$ and $j \in \mathbb{Z}$. Thus (3.1) is true when $n = 0$. This completes the base step.

For the inductive step, fix $n \in \mathbb{Z}^{\geq 0}$. We assume the inductive hypothesis, which is

$$3 \mid (n^3 + 2n). \tag{IH}$$

We will show the inductive conclusion, which is

$$3 \mid ((n+1)^3 + 2(n+1)). \tag{IC}$$

By the (IH), there is a $j \in \mathbb{Z}$ such that

$$n^3 + 2n = 3j. \tag{3.2}$$

From algebra we get

$$\begin{aligned} (n+1)^3 + 2(n+1) &= (n^3 + 3n^2 + 3n + 1) + (2n + 2) \\ &= [n^3 + 2n] + (3n^2 + 3n + 1 + 2) \\ &= [n^3 + 2n] + 3(n^2 + n + 1) \end{aligned}$$

and by (3.2) (which came from the (IH)) and then algebra

$$\begin{aligned} &= [3j] + 3(n^2 + n + 1) \\ &= 3(j + n^2 + n + 1) \\ &= 3k \end{aligned} \tag{3.3}$$

where $k = j + n^2 + n + 1$. Since $k \in \mathbb{Z}$, the calculation through (3.3) show that the inductive conclusion (IC) holds. This completes the inductive step.

Thus the base step and inductive step hold. \square

4. **Theorem 4.** If $n \in \mathbb{Z}$ with $n \geq 0$, then

$$8 \mid (9^n - 1).$$

4.1. Symbolically write Theorem 4. (Do not forget your quantifiers.)

4.2. Prove Theorem 4 using math induction.

hint. See book Proposition 4.4 (p178) and (starred) ER 3.4.8a (p181).

sw. $(\forall n \in \mathbb{Z}^{\geq 0}) [8 \mid (9^n - 1)]$

Proof. We will show that if $n \in \mathbb{Z}^{\geq 0}$, then

$$8 \mid (9^n - 1) \tag{4.1}$$

by basic induction on the variable n .

For the base step, let $n = 0$. Then

$$9^0 - 1 = 9^0 - 1 = 1 - 1 = 0.$$

Also

$$8 \mid 0$$

since $0 = 8j$ where $j = 0$ and $j \in \mathbb{Z}$. Thus (4.1) is true when $n = 0$. This completes the base step.

For the inductive step, fix $n \in \mathbb{Z}^{\geq 0}$. We assume the inductive hypothesis, which is

$$8 \mid (9^n - 1). \tag{IH}$$

We will show the inductive conclusion, which is

$$8 \mid (9^{n+1} - 1). \tag{IC}$$

By the (IH), there is a $j \in \mathbb{Z}$ such that

$$9^n - 1 = 8j$$

and so

$$9^n = 8j + 1. \tag{4.2}$$

From algebra we get

$$9^{n+1} - 1 = 9[9^n] - 1$$

and by (4.2) (which came from the (IH)) and then algebra

$$\begin{aligned} &= 9[8j + 1] - 1 \\ &= (9)(8)j + 9 - 1 \\ &= (8)(9)j + 8 \\ &= 8(9j + 1) \\ &= 8k \end{aligned} \tag{4.3}$$

where $k = 9j + 1$. Since $k \in \mathbb{Z}$, the calculation through (4.3) show the inductive conclusion (IC) holds. This completes the inductive step.

Thus the base step and inductive step hold. \square

5. **Theorem 5.** Let $\{x_n\}_{n=1}^{\infty}$ be the recursively defined sequence defined by

$$x_1 = 1 \quad , \quad x_2 = 2$$

and

$$\text{when } n \in \mathbb{N}, \quad x_{n+2} = \frac{x_{n+1} + x_n}{2}. \tag{RD}$$

If $n \in \mathbb{N}$, then

$$1 \leq x_n \leq 2.$$

5.1. Symbolically write Theorem 5. (Do not forget your quantifiers.)

5.2. Prove Theorem 5 using math induction.

hint. Note you need to prove both the lower bound and upper bound.

$$\text{sw. } \left[\left(x_1 = 1 \wedge x_2 = 2 \wedge (\forall n \in \mathbb{N}) \left[x_{n+2} = \frac{x_{n+1} + x_n}{2} \right] \right) \Rightarrow (\forall n \in \mathbb{N}) [1 \leq x_n \leq 2] \right]$$

Proof. Let the recursively defined sequence $\{x_n\}_{n=0}^\infty$ be given by

$$x_1 = 1 \quad , \quad x_2 = 2$$

and

$$\text{when } n \in \mathbb{N}, \quad x_{n+2} = \frac{x_{n+1} + x_n}{2} . \quad (\text{RD})$$

We will prove that if $n \in \mathbb{N}$ then

$$1 \leq x_n \leq 2 \quad (5.1)$$

by strong induction on n .

For the base step, consider the cases $n = 1$ and $n = 2$. Then you have

$$1 \leq x_1 = 1 \leq 2 \quad \text{and} \quad 1 \leq x_2 = 2 \leq 2.$$

Thus (5.1) holds for $n = 1$ and $n = 2$. This concludes the base step.

For the inductive step, fix $n \in \mathbb{N}^{\geq 2}$. Assume the inductive hypothesis, which is

$$\text{if } j \in \{1, 2, \dots, n\}, \text{ then } 1 \leq x_j \leq 2. \quad (\text{IH})$$

We will show the inductive conclusion, which is

$$1 \leq x_{n+1} \leq 2. \quad (\text{IC})$$

Note $n, n-1 \in \mathbb{N}^{\leq n}$ since $n \geq 2$ and $n \in \mathbb{N}$. Thus by the inductive hypothesis (IH),

$$1 \leq x_n \leq 2 \quad \text{and} \quad 1 \leq x_{n-1} \leq 2. \quad (5.2)$$

The inequalities in (5.2) we gives

$$2 \leq x_n + x_{n-1} \leq 4,$$

which gives

$$1 \leq \frac{x_n + x_{n-1}}{2} \leq 2. \quad (5.3)$$

Since $\frac{x_n + x_{n-1}}{2} = x_{n+1}$, the inequalities in (5.3) give

$$1 \leq x_{n+1} \leq 2, \quad (5.4)$$

which is the inductive conclusion (IC).

By strong induction, we have proven that (5.1) holds for all $n \in \mathbb{N}$. □

6. Theorem 6. Let $\{x_n\}_{n=1}^\infty$ be the recursively defined sequence be

$$x_1 = 1 \quad , \quad x_2 = 1 \quad , \quad x_3 = 1$$

and

$$\text{when } n \in \mathbb{N}^{\geq 4}, \quad x_n = x_{n-1} + x_{n-2} + x_{n-3} . \quad (\text{RD})$$

If $n \in \mathbb{N}$, then

$$x_n < 2^n .$$

6.1. Symbolically write Theorem 6. (Do not forget your quantifiers.)

6.2. Prove Theorem 6 using math induction.

$$\text{sw. } \left[(x_1 = 1 \wedge x_2 = 1 \wedge x_3 = 1 \wedge (\forall n \in \mathbb{N}^{\geq 4}) [x_n = x_{n-1} + x_{n-2} + x_{n-3}]) \Rightarrow (\forall n \in \mathbb{N}) [x_n < 2^n] \right]$$

Proof. Let the recursively defined sequence $\{x_n\}_{n=0}^\infty$ be given by

$$x_1 = 1 \quad , \quad x_2 = 1 \quad , \quad x_3 = 1$$

and

$$\text{when } n \in \mathbb{N}^{\geq 4}, \quad x_n = x_{n-1} + x_{n-2} + x_{n-3} . \quad (\text{RD})$$

We will show that if $n \in \mathbb{N}$, then

$$x_n < 2^n \quad (6.1)$$

by strong induction on n .

For the base step, first let $n = 1$. Then

$$x_n = 1 < 2 = 2^1 = 2^n.$$

When $n = 2$ we have

$$x_n = 1 < 4 = 2^2 = 2^n.$$

When $n = 3$ we have

$$x_n = 1 < 8 = 2^3 = 2^n.$$

Thus (6.1) holds for $n = 1$ and $n = 2$ and $n = 3$. This concludes the base step.

For the inductive step, fix $n \in \mathbb{N}^{\geq 3}$. Assume the inductive hypothesis, which is

$$\text{if } j \in \{1, 2, \dots, n\}, \text{ then } x_j < 2^j. \quad (\text{IH})$$

We will show the inductive conclusion, which is

$$x_{n+1} < 2^{n+1}. \quad (\text{IC})$$

The recursive definition (RD) is valid for x_k when $k \geq 4$. We have $n \geq 3$, so $n + 1 \geq 4$. So by (RD)

$$x_{n+1} = x_n + x_{n-1} + x_{n-2}$$

Since $3 \leq n$ we have $1 \leq n - 2$ so $1 \leq n - 2 \leq n - 1 \leq n$. Thus we can apply the inductive hypothesis (IH) to n , $n - 1$ and $n - 2$ since

$$n, n - 1, n - 2 \in \{1, 2, \dots, n\}.$$

Thus we have by the recursive definition

$$x_{n+1} = x_n + x_{n-1} + x_{n-2}$$

and by (IH)

$$< 2^n + 2^{n-1} + 2^{n-2}$$

and by algebra (look at where you are going: want a 2^{n+1})

$$= 2^{n+1} (2^{-1} + 2^{-2} + 2^{-3})$$

$$= 2^{n+1} \left(\frac{1}{2} + \frac{1}{4} + \frac{1}{8} \right)$$

$$= 2^{n+1} \left(\frac{4 + 2 + 1}{8} \right)$$

$$= 2^{n+1} \left(\frac{7}{8} \right)$$

$$< 2^{n+1}. \quad (6.2)$$

The calculations through (6.2) show the inductive conclusion (IC) holds. This completes the inductive step.

Thus the base step and the inductive step hold. \square

Remark. The calculation in (6.2) could also have been done as below.

$$x_{n+1} = x_n + x_{n-1} + x_{n-2}$$

$$< 2^n + 2^{n-1} + 2^{n-2}$$

$$= 2^{n-2} (2^2 + 2^1 + 2^0)$$

$$= 2^{n-2} (7)$$

$$< 2^{n-2} (2^3)$$

$$= 2^{n+1}.$$