

# Lecture 5: Classifying Reducible Polynomials with Small Norm

**Theorem (Schinzel):** Fix  $a_0, \dots, a_r \in \mathbb{Z} - \{0\}$ . Then there is an algorithm for obtaining a finite classification of the polynomials of the form  $a_r x^{d_r} + \dots + a_1 x^{d_1} + a_0$  that have reducible non-reciprocal part.

**Lemma:** Let  $s$  and  $t$  be positive integers. Suppose a system of linear equations in the variables  $x_0, \dots, x_s$  is of the form

$$\alpha_{i0}x_0 + \alpha_{i1}x_1 + \dots + \alpha_{is}x_s = \beta_i \quad \text{for } 1 \leq i \leq t,$$

where the  $\alpha_{ij}$  and  $\beta_i$  are all in  $\mathbb{Z}$ . Suppose further that the system of equations has infinitely many solutions  $(x_0, \dots, x_s) \in \mathbb{R}^{s+1}$ . If the system has at least one solution  $(x_0, \dots, x_s) \in \mathbb{Z}^{s+1}$  with  $x_0, x_1, \dots, x_s$  *distinct*, then the system has infinitely many such solutions.

**Main Ideas for Proof of Lemma:**

- Set  $A = (\alpha_{i,j-1})$ , a  $t \times (s+1)$  matrix, and let  $\rho$  be its rank.
- Rearrange so the the first  $\rho$  rows and first  $\rho$  columns are linearly independent.
- Let  $B$  be the  $\rho \times \rho$  matrix from the upper left part of  $A$ , and note that  $D = |\det B| \geq 1$ .
- Solve to obtain  $x_i = \frac{1}{D} \left( c_i + \sum_{j=\rho}^s b_{ij} x_j \right)$  for  $0 \leq i \leq \rho - 1$  with  $c_i$  and  $b_{ij}$  in  $\mathbb{Z}$ .
- Fix a solution  $(k_0, k_1, \dots, k_s)$  consisting of distinct integers.
- Define  $k'_i = k_i + \ell_i D$  for  $\rho \leq i \leq s$ , and  $k'_i = \frac{1}{D} \left( c_i + \sum_{j=\rho}^s b_{ij} k'_j \right) = k_i + \sum_{j=\rho}^s b_{ij} \ell_j$  for  $0 \leq i \leq \rho - 1$ . Note that  $(k'_0, k'_1, \dots, k'_s)$  is a solution and each  $k_j \in \mathbb{Z}$ .
- Prove the  $k'_j$ 's are distinct by taking  $\ell_j \equiv 0 \pmod{d}$ , for all  $j$ , where  $d$  is large (so that the  $k'_j$ 's are distinct modulo  $d$ ).

**Proof of Theorem:**

- First, consider the case that the  $d_j$  (and  $a_j$ ) are fixed.
- Recall the non-reciprocal part of  $f(x)$  is reducible if and only if there exists  $w(x)$  different from  $\pm f(x)$  and  $\pm \tilde{f}(x)$  such that  $w(x)\tilde{w}(x) = f(x)\tilde{f}(x)$ .
- Write  $f(x) = \sum_{j=0}^r a_j x^{d_j}$  and  $w(x) = \sum_{j=0}^s b_j x^{k_j}$ . Here, the  $a_j$  and  $d_j$  are given integers with  $0 = d_0 < d_1 < \dots < d_{r-1} < d_r = n$ ; the  $b_j$  and  $k_j$  as unknown integers with  $0 = k_0 < k_1 < \dots < k_{s-1} < k_s = n$ .

- Since  $\|w\| = \|f\|$ , we deduce  $\sum_{j=0}^s |b_j| \leq \|f\|^2$ . Thus, there are finitely many possibilities for the  $b_j$ 's. Fix the  $b_j$ 's.
- Define  $E = \{n - k_j + k_i : 0 \leq i, j \leq s\}$ , the set of exponents appearing in  $w(x)\tilde{w}(x)$ . Consider a system of equations with each equation consisting of an element from  $E$  equal to either another element of  $E$  (possible cancellation) or an element of  $E$  equal to an expression of the form  $n - d_j + d_i$  (from the right-hand side of  $w(x)\tilde{w}(x) = f(x)\tilde{f}(x)$ ). Consider only a system satisfying: (i) each element of  $E$  occurs in such an equation at least once, (ii) every exponent of an uncanceled term in  $f(x)\tilde{f}(x)$  is used exactly once, and (iii) the equations  $n - k_s + k_0 = 0$  and  $n - k_0 + k_s = 2n$  are used. We only allow equations of the form  $n - k_j + k_i = n - k_v + k_u$  if  $(i, j) \neq (u, v)$ . Replace the equations in (iii) with  $k_0 = 0$  and  $k_s = n$ . We want to know if the system has a solution (for each such system).
- One of the following three possibilities for a system may occur: (i') the system may have a unique solution (in  $\mathbb{R}^{s+1}$ ), (ii') the system may have no solutions, or (iii') the system may have infinitely many solutions. The cases (i') and (ii') are good.
- Justify (iii') is impossible in distinct integers  $k_j$ . By the lemma, there is a solution in distinct integers  $k'_j$  with either  $k'_u = \min_{0 \leq j \leq s} \{k'_j\} \leq -1$  or  $k'_v = \max_{0 \leq j \leq s} \{k'_j\} \geq n + 1$ . Note both  $k'_u \leq 0$  and  $k'_v \geq n$  hold. Hence,  $n - k'_v + k'_u \leq -1$ . Either  $n - k'_v + k'_u = n - k'_j + k'_i$  with  $(i, j) \neq (u, v)$  or  $n - k'_v + k'_u = m$  for some exponent  $m$  appearing in  $f(x)\tilde{f}(x)$ . Both are impossibilities.
- For variable  $d_j$ , consider each possibility of cancelled terms in  $f(x)\tilde{f}(x)$  and proceed as above.
- After obtaining a solution for  $w(x)$ , plug the result into  $w(x)\tilde{w}(x) = f(x)\tilde{f}(x)$  and solve for the  $d_j$ . Here, the possibility of infinitely many solutions in the  $d_j$  is fine (and occurs).
- Plug in the resulting  $d_j$  to see if now  $w(x) = \pm f(x)$  or  $w(x) = \pm \tilde{f}(x)$ . This requires solving another system of equations. Discuss what the final classification looks like.