



# Probabilistic Methods in Combinatorics Lecture 7

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- Dependence graph:  $d_{SS} \leq 3n$ ,  $d_{ST} \leq 3\binom{n}{k-2}$ ,  $d_{TS} \leq \binom{k}{2}n$ , and  $d_{TT} \leq \binom{k}{2}\binom{n}{k-2}$ .



# Proof

By LLL, we only require

$$\begin{aligned} p^3 &\leq x(1-x)^{3n}(1-y)^{3\binom{n}{k-2}} \\ (1-p)^{\binom{k}{2}} &\leq y(1-x)^{\binom{k}{2}n}(1-y)^{\binom{k}{2}\binom{n}{k-2}}. \end{aligned}$$



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We can choose  $p = c_1 n^{-1/2}$ ,  $k = c_2 n^{1/2} \log n$ ,  $x = c_3 n^{-3/2}$ ,  
and  $y = c_4 / \binom{n}{k}$ .





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This gives  $R(3, k) > c_5 k^2 / \log^2 k$ . □



# $R(4, k)$

Best bounds for  $R(r, k)$  (for fixed  $r$  and  $k$  large),

$$c \left( \frac{k}{\log k} \right)^{(r+1)/2} < R(r, k) < (1 + o(1)) \frac{k^{r-1}}{\log^{r-2} k}.$$

**Erdős conjecture \$250:** Prove

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The best lower bound is using LLL;  $R(4, k) > c' \frac{k^{2.5}}{\log^{2.5} k}$ .



# Directed cycles

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**Proof:** First we can assume every out-degree is  $\delta$  by deleting some edges if necessary. Consider  $f: V \rightarrow \mathbb{Z}_k$ . Bad event  $A_v$ : no  $u \in \Gamma^+(v)$  with  $f(u) = f(v) + 1$ .

$$\Pr(A_v) = (1 - 1/k)^\delta.$$



Each event depends on at most  $\delta\Delta$  others. Apply LLL.  $\square$

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If the conjecture is true, then it is tight.

$$la(G) \geq \frac{nd}{2(n-1)} > \frac{d}{2}.$$



# Directed graphs

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DLA conjecture for  $d$  implies LA conjecture for  $2d$ .



# A proposition

**Proposition:** *Let  $H = (V, E)$  be a graph with maximum degree  $d$ , and let  $V = V_1 \cup V_2 \cup \dots \cup V_r$  be a partition of  $V$ . If  $|V_i| \geq 2ed$ , then there is an independent set of vertices  $W$  that contains a vertex from each  $V_i$ .*



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**Proof:** WLOG, we assume

$$|V_1| = |V_2| = \dots = |V_r| = \lceil 2ed \rceil = g.$$

Pick from each  $V_i$  a vertex randomly and independently. Let  $W$  be the random set of the vertices picked. For each edge  $f$ , let  $A_f$  be the event that both ends in  $W$ . The maximum degree in the dependence graph is at most  $2gd - 1$ . We have  $e \cdot 2gd \cdot \frac{1}{g^2} = \frac{2ed}{g} < 1$ . Apply LLL.  $\square$



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**Proof:** Using Hall's matching theorem, we can partition  $F$  into  $d$  pairwise disjoint 1-regular spanning subgraphs  $F_1, \dots, F_d$  of  $G$ .



# Continue

Each  $F_i$  is a union of vertex disjoint directed cycles. Let  $V_1, \dots, V_r$  are the sets of edges of all cycles. Then

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Apply the proposition to the line-graph  $H$  of  $G$ . Note  $H$  is  $4d - 2$ -regular.



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There exists an independent set  $M_1$  of  $H$ . Now  $M_1, F_1 \setminus M_1, \dots, F_d \setminus M_1$  forms  $d + 1$  linear directed forests.

□



# General $d$ -regular graphs

**Theorem [Alon 1988]** There is an absolute constant  $c > 0$  such that for every  $d$ -regular directed graph  $G$

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**Corollary** There is an absolute constant  $c > 0$  such that for every  $d$ -regular graph  $G$

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The error terms can be improved to  $cd^{2/3} \log^{1/3} d$ .





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Define for  $i \in \mathbb{Z}_p$ ,

$$E_i = \{(u, v) \in E : f(v) = f(u) + i\}.$$

Let  $G_i = (V, E_i)$  and

- $\Delta_i^+$ : the maximum out-degree of  $G_i$ .
- $\Delta_i^-$ : the maximum in-degree of  $G_i$ .
- $\Delta_i$ : the maximum of  $\Delta_i^+$  and  $\Delta_i^-$ .



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There exists a  $f$  satisfying

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$$\text{dla}(G) \leq 2\Delta_0 + \sum_{i=1}^{p-1} (\Delta_i + 1) \leq d + d/p + p + C\sqrt{dp \log d}.$$

Now choose  $p \sim d^{1/2}$ .



# Chernoff Inequality

Suppose  $X = \sum_{i=1}^n X_i$ , where  $X_i$  are independent 0-1 random variables. Then we have

$$\Pr(X < E(X) - \lambda) \leq e^{-\frac{\lambda^2}{2E(X)}}$$
$$\Pr(X > E(X) + \lambda) \leq e^{-\frac{\lambda^2}{2(E(X) + \lambda/3)}}$$

