

Definitions

Def. Two statements \tilde{P} and \tilde{Q} are **logically equivalent** provided they have the same truth value for all possible combinations of truth values for all variables (i.e. all atoms) appearing in \tilde{P} and \tilde{Q} . In this case, we write $\tilde{P} \equiv \tilde{Q}$ and say that \tilde{P} and \tilde{Q} are logically equivalent. p43

Def. • The **converse** of the conditional statement $P \Rightarrow Q$ is the conditional statement $Q \Rightarrow P$. p44
 • The **contrapositive** of the conditional statement $P \Rightarrow Q$ is the condition statement $(\sim Q) \Rightarrow (\sim P)$.
 • **Rmk.** We have already seen that: $[P \Rightarrow Q] \not\equiv [Q \Rightarrow P]$ but $[P \Rightarrow Q] \equiv [(\sim Q) \Rightarrow (\sim P)]$.

Def. A **negation** (also called **denial**) of a statement P is $\sim P$. p33

Recall. The priority of connectives is: \sim (high) , \wedge , \vee , \Rightarrow , \Leftrightarrow (low). So $\sim P \vee \sim Q$ is $(\sim P) \vee (\sim Q)$.

Important Logical Equivalencies

Theorem 2.8. Let P , Q , and R be statements. Thm 2.8

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Double Negation:

$$[\sim(\sim P)] \equiv P. \quad (1)$$

Biconditional Statement:

$$[P \Leftrightarrow Q] \equiv [(P \Rightarrow Q) \wedge (Q \Rightarrow P)]. \quad (2)$$

De Morgans Laws:

$$[\sim(P \wedge Q)] \equiv [(\sim P) \vee (\sim Q)] \quad (3)$$

$$[\sim(P \vee Q)] \equiv [(\sim P) \wedge (\sim Q)]. \quad (4)$$

Distributive Laws:

$$[P \vee (Q \wedge R)] \equiv [(P \vee Q) \wedge (P \vee R)] \quad (5)$$

$$[P \wedge (Q \vee R)] \equiv [(P \wedge Q) \vee (P \wedge R)]. \quad (6)$$

Conditional Statements:

$$[P \Rightarrow Q] \equiv [(\sim Q) \Rightarrow (\sim P)] \quad (\text{contrapositive}) \quad (7)$$

$$[P \Rightarrow Q] \equiv [(\sim P) \vee Q] \quad (\text{how do you keep a promise?}) \quad (8)$$

$$[\sim(P \Rightarrow Q)] \equiv [P \wedge (\sim Q)] \quad (\text{how do you break a promise?}) \quad (9)$$

$$[\sim(P \wedge Q)] \equiv [P \Rightarrow (\sim Q)]. \quad (\text{not in book}) \quad (10)$$

Conditionals with Disjunctions:

$$[(P \vee Q) \Rightarrow R] \equiv [(P \Rightarrow R) \wedge (Q \Rightarrow R)] \quad (11)$$

$$[P \Rightarrow (Q \vee R)] \equiv [(P \wedge (\sim Q)) \Rightarrow R]. \quad (12)$$