

1. THE HOMOTOPY CATEGORY

Corollary 1. *Let \mathcal{C} be a model category. Let $\gamma: \mathcal{C}_{cf} \rightarrow \text{Ho}(\mathcal{C}_{cf})$ and $\delta: \mathcal{C}_{cf} \rightarrow \mathcal{C}_{cf}/\sim$ be the canonical functors. Then there is a unique isomorphism of categories making the diagram*

$$\begin{array}{ccc}
 \mathcal{C}_{cf} & \xrightarrow{\gamma} & \text{Ho}(\mathcal{C}_{cf}) \\
 & \searrow \delta & \nearrow \exists! j \\
 & & \mathcal{C}_{cf}/\sim
 \end{array}$$

commute. Furthermore, j is the identity on objects.

Proof. We show that \mathcal{C}_{cf}/\sim satisfies the same universal property as $\text{Ho}(\mathcal{C}_{cf})$. We note that δ takes homotopy equivalences to isomorphisms and so by Proposition 1.2.8, δ also takes weak equivalences to isomorphisms. Let $\mathcal{F}: \mathcal{C}_{cf} \rightarrow \mathcal{D}$ be a functor that takes weak equivalences to isomorphisms. Given an object X , take the functorial cylinder object $X \times I$. This guarantees that $X \times I$ is cofibrant and fibrant, since $X \amalg X$ is also cofibrant and X is fibrant. Now note that we have a commutative diagram

$$\begin{array}{ccccc}
 & & \text{id}_X & & \\
 & & \curvearrowright & & \\
 X & \begin{array}{c} \xrightarrow{u_0} \\ \xrightarrow{u_1} \end{array} & X \amalg X & \xrightarrow{\nabla} & X \\
 & \searrow & \searrow \text{\scriptsize } \nu_0 \amalg \nu_1 & \searrow s & \\
 & & X \times I & & \\
 & \swarrow & \swarrow \text{\scriptsize } \iota_0 & \swarrow & \\
 & & & &
 \end{array}$$

Since s and id_X are weak equivalences we have that ι_0 and ι_1 are both weak equivalences by 2-out-of-3. Hence

$$\mathcal{F}(s) \circ \mathcal{F}(\iota_0) = \mathcal{F}(s \circ \iota_0) = \text{id}_{\mathcal{F}X} = \mathcal{F}(s \circ \iota_1) = \mathcal{F}(s) \circ \mathcal{F}(\iota_1)$$

implies $\mathcal{F}(\iota_0) = \mathcal{F}(\iota_1)$.

Given a homotopy

$$\begin{array}{ccc}
 X & \begin{array}{c} \xrightarrow{\iota_0} \\ \xrightarrow{\iota_1} \end{array} & X \times I \\
 & \searrow \text{\scriptsize } f_0 & \downarrow H \\
 & & Y
 \end{array}$$

we have

$$\mathcal{F}(f_0) = \mathcal{F}(H \circ \iota_0) = \mathcal{F}(H) \circ \mathcal{F}(\iota_0) = \mathcal{F}(H) \circ \mathcal{F}(\iota_1) = \mathcal{F}(H \circ \iota_1) = \mathcal{F}(f_1)$$

so that \mathcal{F} identifies left homotopic maps. By duality, we see that \mathcal{F} also identifies right homotopic maps. This allows us to define a unique functor $\mathcal{G}: \mathcal{C}_{cf}/\sim \rightarrow \mathcal{D}$ by $\mathcal{G}\delta X = \mathcal{F}X$ and $\mathcal{G}\delta(f) = \mathcal{F}(f)$. The result now follows by unicity of universal objects. ■

Theorem 1. *Suppose \mathcal{C} is a model category. Let $\gamma: \mathcal{C} \rightarrow \text{Ho}(\mathcal{C})$ denote the canonical functor, Q the cofibrant replacement functor of \mathcal{C} , and R the fibrant replacement functor.*

(i) *The inclusion $\mathcal{C}_{cf} \rightarrow \mathcal{C}$ induces an equivalence of categories*

$$\mathcal{C}_{cf}/\sim \xrightarrow{\sim} \text{Ho}(\mathcal{C}_{cf}) \longrightarrow \text{Ho}(\mathcal{C})$$

(ii) There are natural isomorphisms

$$\mathcal{C}(QRX, QRY)/\sim \xrightarrow{\sim} \mathrm{Ho}(\mathcal{C})(\gamma X, \gamma Y) \xrightarrow{\sim} \mathcal{C}(RQX, RQY)/\sim$$

In addition, there is a natural isomorphism $\mathrm{Ho}(\mathcal{C})(\gamma X, \gamma Y) \cong \mathcal{C}(QX, RY)/\sim$, and, if X is cofibrant and Y is fibrant, there is a natural isomorphism $\mathrm{Ho}(\mathcal{C})(\gamma X, \gamma Y) \cong \mathcal{C}(X, Y)/\sim$. In particular, $\mathrm{Ho}(\mathcal{C})$ is a category without moving to a higher universe.

(iii) The functor $\gamma: \mathcal{C} \rightarrow \mathrm{Ho}(\mathcal{C})$ identifies left or right homotopic maps.

(iv) If $f: A \rightarrow B$ is a map in \mathcal{C} such that γf is an isomorphism in $\mathrm{Ho}(\mathcal{C})$, then f is a weak equivalence.

Proof. Part (i) is just the composition of the isomorphism from the Corollary above and the equivalence $\mathrm{Ho}(\mathcal{C}_{cf}) \rightarrow \mathrm{Ho}(\mathcal{C})$ of Proposition 1.2.3.

The first part of (ii) follows from the equivalences

$$\mathrm{Ho}(\mathcal{C}) \xrightarrow{\mathrm{Ho}(Q)} \mathrm{Ho}(\mathcal{C}_c) \xrightarrow{\mathrm{Ho}(R)} \mathrm{Ho}(\mathcal{C}_{cf}) \quad \text{and} \quad \mathrm{Ho}(\mathcal{C}) \xrightarrow{\mathrm{Ho}(R)} \mathrm{Ho}(\mathcal{C}_f) \xrightarrow{\mathrm{Ho}(Q)} \mathrm{Ho}(\mathcal{C}_{cf})$$

The rest follows from Proposition 1.2.5 and the fact that

$$QX \xrightarrow{q_X} X \xrightarrow{r_X} RX$$

is a weak equivalence.

Part (iii) was proved in Corollary 1.2.9.

For part (iv), assume that $f \in \mathcal{C}(X, Y)$ is a morphism such that γf is an isomorphism in $\mathrm{Ho}(\mathcal{C})$. This implies that the image of QRf in \mathcal{C}_{cf}/\sim is an isomorphism, and from Corollary 1.2.7 it follows that QRf is a homotopy equivalence. By Proposition 1.2.8, QRf is a weak equivalence and so the left-hand side of the commutative diagram

$$\begin{array}{ccccc} QRX & \xrightarrow{q_{RX}} & RX & \xleftarrow{r_X} & X \\ \downarrow QRf & & \downarrow Rf & & \downarrow f \\ QRY & \xrightarrow{q_{RY}} & RY & \xleftarrow{r_Y} & Y \end{array}$$

implies by 2-out-of-3 that Rf is a weak equivalence, and thus the right-hand square implies by 2-out-of-3 that f is a weak equivalence, as desired. \blacksquare

2. ADJOINTS

Lemma 1. Let $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$ be a functor and let $\mathcal{D}_0: \mathcal{D}_0 \rightarrow \mathcal{D}$ be a full subcategory of \mathcal{D} . If for each object X of \mathcal{C} there exists an object Y of \mathcal{D}_0 and an isomorphism $\mathcal{F}X \xrightarrow{\sim} \mathcal{D}_0(Y)$, then there exists a functor $\mathcal{F}_0: \mathcal{C} \rightarrow \mathcal{D}_0$ and a natural isomorphism $\theta_0: \mathcal{F} \rightarrow \mathcal{D}_0 \circ \mathcal{F}_0$. Moreover, \mathcal{F}_0 is unique up to unique isomorphism.

Proof. For each object X of \mathcal{C} , choose an object Y of \mathcal{D}_0 such that $\theta_0(X): \mathcal{F}X \cong \mathcal{D}_0(Y)$ and define $\mathcal{F}_0(X) = Y$. Since \mathcal{F} is a functor we have a morphism

$$\mathcal{F}(X_1, X_2): \mathcal{C}(X_1, X_2) \rightarrow \mathcal{D}(\mathcal{F}X_1, \mathcal{F}X_2).$$

for each pair of objects X_1, X_2 of \mathcal{C} , and since \mathcal{D}_0 is a full subcategory, we also have an isomorphism

$$\mathcal{D}_0(Y_1, Y_2): \mathcal{D}_0(Y_1, Y_2) \xrightarrow{\sim} \mathcal{D}(\mathcal{D}_0 Y_1, \mathcal{D}_0 Y_2)$$

for any pair of objects Y_1, Y_2 of \mathcal{D}_0 . We thus obtain a morphism

$$\begin{array}{ccc}
\mathcal{C}(X_1, X_2) & \xrightarrow{\mathcal{F}(X_1, X_2)} & \mathcal{D}(\mathcal{F}X_1, \mathcal{F}X_2) \\
& & \downarrow h_{\mathcal{F}X_2}(\theta_0(X_1)^{-1}) \\
& \begin{array}{ccc} f \longmapsto & \mathcal{F}(f) & \\ & \downarrow & \\ & \mathcal{F}(f) \circ \theta_0(X_1)^{-1} & \\ & \downarrow & \\ & \theta_0(X_2) \circ \mathcal{F}(f) \circ \theta_0(X_1)^{-1} & \\ & \downarrow & \\ & \mathcal{F}_0(X_1, X_2) \circ \mathcal{F}_0^{-1}(\theta_0(X_2) \circ \mathcal{F}(f) \circ \theta_0(X_1)^{-1}) & \end{array} & \begin{array}{ccc} \mathcal{D}(\mathcal{F}X_1, \mathcal{F}X_2) & & \\ & \downarrow h^{i_0 \circ \mathcal{F}_0 X}(\theta_0(X_2)) & \\ & \mathcal{D}(i_0 \circ \mathcal{F}_0 X_1, i_0 \circ \mathcal{F}_0 X_2) & \\ & \downarrow i_0(\mathcal{F}_0 X_1, \mathcal{F}_0 X_2)^{-1} & \\ & \mathcal{D}_0(\mathcal{F}_0 X_1, \mathcal{F}_0 X_2) & \end{array} \\
& \searrow \mathcal{F}_0(X_1, X_2) & \rightarrow
\end{array}$$

Since all the morphisms involved respect composition, it's clear that this assignment is functorial and by construction makes the isomorphism $\theta_0: \mathcal{F} \rightarrow i_0 \circ \mathcal{F}_0$ natural.

For unicity, suppose that $\mathcal{G}_0: \mathcal{C} \rightarrow \mathcal{D}_0$ is another functor equipped with a natural isomorphism $\varphi_0: \mathcal{F} \rightarrow i_0 \circ \mathcal{G}_0$. We have by assumption an isomorphism of functors

$$i_0 \circ G_0 \xrightarrow{\varphi_0^{-1}} \mathcal{F} \xrightarrow{\theta_0} i_0 \mathcal{F}_0.$$

Since i_0 is fully faithful, for each object X of \mathcal{C} it reflects the isomorphism $\theta_0(X) \circ \varphi_0(X)^{-1}$ giving an isomorphism

$$i_0^{-1}(\theta_0(X) \circ \varphi_0(X)^{-1}) = \eta(X): \mathcal{G}_0(X) \longrightarrow \mathcal{F}_0(X).$$

To see that this is natural, we consider the diagram

$$\begin{array}{ccc}
X & & \mathcal{G}_0(X) \xrightarrow{\eta(X)} \mathcal{F}_0(X) \\
\downarrow f & & \downarrow \mathcal{G}_0(f) \quad \downarrow \mathcal{F}_0(f) \\
X' & & \mathcal{G}_0(X') \xrightarrow{\eta(X')} \mathcal{F}_0(X')
\end{array}$$

and note that from the two naturality squares

$$\begin{array}{ccc}
\mathcal{F}(X) \xrightarrow{\theta_0(X)} i_0 \circ \mathcal{F}_0(X) & & \mathcal{F}(X) \xrightarrow{\varphi_0(X)} i_0 \circ \mathcal{G}_0(X) \\
\downarrow \mathcal{F}(f) & \downarrow i_0 \circ \mathcal{F}_0(f) & \downarrow \mathcal{F}(f) & \downarrow i_0 \circ \mathcal{G}_0(f) \\
\mathcal{F}(X') \xrightarrow{\theta_0(X')} i_0 \circ \mathcal{F}_0(X') & & \mathcal{F}(X') \xrightarrow{\varphi_0(X')} i_0 \circ \mathcal{G}_0(X')
\end{array} \text{ and }$$

we obtain

$$\begin{aligned}
\iota_0(\mathcal{F}_0(f) \circ \eta(X)) &= \iota_0 \circ \mathcal{F}_0(f) \circ \theta_0(X) \circ \varphi_0(X)^{-1} \\
&= \theta_0(X') \circ \mathcal{F}(f) \circ \theta_0(X)^{-1} \circ \theta_0(X) \circ \varphi_0(X)^{-1} \\
&= \theta_0(X') \circ \mathcal{F}(f) \circ \varphi_0(X)^{-1} \\
&= \theta_0(X') \circ \varphi_0(X')^{-1} \circ \iota_0 \circ \mathcal{G}_0(f) \\
&= \iota_0(\eta(X')) \circ \iota_0 \circ \mathcal{G}_0(f) \\
&= \iota_0(\eta(X') \circ \mathcal{G}_0(f))
\end{aligned}$$

which implies $\mathcal{F}_0(f) \circ \eta(X) = \eta(X') \circ \mathcal{G}_0(f)$ because ι_0 is fully faithful. ■

Proposition 1. Consider the two functors $\mathcal{L}: \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$. The following are equivalent:

(i) for each object X of \mathcal{C} and Y of \mathcal{D} , there is an isomorphism

$$\mathcal{D}(\mathcal{L}X, Y) \cong \mathcal{C}(X, \mathcal{R}Y)$$

natural in both X and Y ,

(ii) there exist natural transformations

$$\varepsilon: \text{id}_{\mathcal{C}} \rightarrow \mathcal{R} \circ \mathcal{L} \text{ and } \eta: \mathcal{L} \circ \mathcal{R} \rightarrow \text{id}_{\mathcal{D}}$$

called the unit and counit of adjunction, respectively, that make the diagrams

$$\begin{array}{ccc}
\mathcal{L}X & \xrightarrow{\mathcal{L}(\varepsilon_X)} & \mathcal{L} \circ \mathcal{R} \circ \mathcal{L}X \\
& \searrow \text{id}_{\mathcal{L}X} & \downarrow \eta_{\mathcal{L}X} \\
& & \mathcal{L}X
\end{array}
\quad \text{and} \quad
\begin{array}{ccc}
\mathcal{R}Y & \xrightarrow{\varepsilon_{\mathcal{R}Y}} & \mathcal{R} \circ \mathcal{L} \circ \mathcal{R}Y \\
& \searrow \text{id}_{\mathcal{R}Y} & \downarrow \mathcal{R}(\eta_Y) \\
& & \mathcal{R}Y
\end{array}$$

commute for all objects X of \mathcal{C} and Y of \mathcal{D} ,

(iii) for every object Y of \mathcal{D} , the functor

$$\mathcal{D}(\mathcal{L}(-), Y): \mathcal{C} \rightarrow \mathbf{Set}$$

is representable by $\mathcal{R}Y$.

If any of these conditions holds, we say that \mathcal{L} is left adjoint to \mathcal{R} (symmetrically, \mathcal{R} is right adjoint to \mathcal{L}), and write $\mathcal{L} \dashv \mathcal{R}$.

Note that condition (iii) implies that the right adjoint of \mathcal{L} is unique up to unique isomorphism and, symmetrically, the left adjoint of \mathcal{R} is as well.

Proof. (i \Rightarrow ii)

For each object X of \mathcal{C} and Y of \mathcal{D} define ε_X and η_Y to be the images of the identity morphism under the isomorphisms

$$\mathcal{D}(\mathcal{L}X, \mathcal{L}X) \xrightarrow{\sim} \mathcal{C}(X, \mathcal{R} \circ \mathcal{L}X) \text{ and } \mathcal{C}(\mathcal{R}Y, \mathcal{R}Y) \xrightarrow{\sim} \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y),$$

respectively. One sees that these are natural transformations by chasing the identities through the commutative diagrams

$$\begin{array}{ccc}
X & & \mathcal{D}(\mathcal{L}X, \mathcal{L}X) \xrightarrow{\sim} \mathcal{C}(X, \mathcal{R} \circ \mathcal{L}X) \\
\downarrow f & & \downarrow h^{\mathcal{L}X \circ \mathcal{L}(f)} \qquad \qquad \downarrow h^X(\mathcal{R} \circ \mathcal{L}(f)) \\
& & \mathcal{D}(\mathcal{L}X, \mathcal{L}X') \xrightarrow{\sim} \mathcal{C}(X, \mathcal{R} \circ \mathcal{L}X') \\
& & \uparrow h_{\mathcal{L}X'}(\mathcal{L}f) \qquad \qquad \uparrow h_{\mathcal{R} \circ \mathcal{L}X'}(f) \\
X' & & \mathcal{D}(\mathcal{L}X', \mathcal{L}X') \xrightarrow{\sim} \mathcal{C}(X', \mathcal{R} \circ \mathcal{L}X')
\end{array}$$

and

$$\begin{array}{ccc}
Y & & \mathcal{C}(\mathcal{R}Y, \mathcal{R}Y) \xrightarrow{\sim} \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y) \\
\downarrow g & & \downarrow h^{\mathcal{R}Y}(\mathcal{R}(g)) \qquad \qquad \downarrow h^{\mathcal{L} \circ \mathcal{R}Y}(g) \\
& & \mathcal{C}(\mathcal{R}Y, \mathcal{R}Y') \xrightarrow{\sim} \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y') \\
& & \uparrow h_{\mathcal{R}Y'}(\mathcal{R}g) \qquad \qquad \uparrow h_{Y'}(\mathcal{L} \circ \mathcal{R}(g)) \\
Y' & & \mathcal{C}(\mathcal{R}Y', \mathcal{R}Y') \xrightarrow{\sim} \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y', Y')
\end{array}$$

To establish the commutativity conditions, we chase the identities through the commutative diagrams

$$\begin{array}{ccc}
X & & \mathcal{C}(\mathcal{R} \circ \mathcal{L}X, \mathcal{R} \circ \mathcal{L}X) \xrightarrow{\sim} \mathcal{D}(\mathcal{L} \circ \mathcal{R} \circ \mathcal{L}X, \mathcal{L}X) \\
\downarrow \varepsilon_X & & \downarrow h_{\mathcal{R} \circ \mathcal{L}X}(\varepsilon_X) \qquad \qquad \downarrow h_{\mathcal{L}X}(\mathcal{L}(\varepsilon_X)) \\
\mathcal{R} \circ \mathcal{L}X & & \mathcal{C}(X, \mathcal{R} \circ \mathcal{L}X) \xrightarrow{\sim} \mathcal{D}(\mathcal{L}X, \mathcal{L}X)
\end{array}$$

and

$$\begin{array}{ccc}
\mathcal{L} \circ \mathcal{R}Y & & \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, \mathcal{L} \circ \mathcal{R}Y) \xrightarrow{\sim} \mathcal{C}(\mathcal{L} \circ \mathcal{L} \circ \mathcal{R}Y, \mathcal{L}Y) \\
\downarrow \eta_Y & & \downarrow h^{\mathcal{L} \circ \mathcal{R}Y}(\eta_Y) \qquad \qquad \downarrow h^{\mathcal{R}Y}(\mathcal{R}(\eta_Y)) \\
Y & & \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y) \xrightarrow{\sim} \mathcal{C}(\mathcal{R}Y, \mathcal{R}Y)
\end{array}$$

and use the definition of ε and η . This establishes (ii).

(ii \Rightarrow iii)

Define the morphisms

$$\mathcal{D}(\mathcal{L}X, Y) \xrightarrow{\mathcal{R}(\mathcal{L}X, Y)} \mathcal{C}(\mathcal{R} \circ \mathcal{L}X, \mathcal{R}Y) \xrightarrow{h_{\mathcal{R}Y}(\varepsilon_X)} \mathcal{C}(X, \mathcal{R}Y)$$

and

$$\mathcal{C}(X, \mathcal{R}Y) \xrightarrow{\mathcal{L}(X, \mathcal{R}Y)} \mathcal{D}(\mathcal{L}X, \mathcal{L} \circ \mathcal{R}Y) \xrightarrow{h^{\mathcal{L}X}(\eta_Y)} \mathcal{D}(\mathcal{L}X, Y).$$

Using the naturality diagrams

$$\begin{array}{ccc}
 X & \xrightarrow{\varepsilon_X} & \mathcal{R} \circ \mathcal{L}X \\
 \downarrow f & & \downarrow \mathcal{R} \circ \mathcal{L}(f) \\
 \mathcal{R}Y & \xrightarrow{\varepsilon_{\mathcal{R}Y}} & \mathcal{R} \circ \mathcal{L} \circ \mathcal{R}Y
 \end{array}
 \quad \text{and} \quad
 \begin{array}{ccc}
 \mathcal{L} \circ \mathcal{R} \circ \mathcal{L}X & \xrightarrow{\eta_{\mathcal{L}X}} & \mathcal{L}X \\
 \downarrow \mathcal{R} \circ \mathcal{L}(g) & & \downarrow g \\
 \mathcal{L} \circ \mathcal{R}Y & \xrightarrow{\eta_Y} & Y
 \end{array}$$

we see

$$\begin{aligned}
 h_{\mathcal{R}Y}(\varepsilon_X) \circ \mathcal{R}(\mathcal{L}X, Y) \circ h^{\mathcal{L}X}(\eta_Y) \circ \mathcal{L}(X, \mathcal{R}Y)(f) &= \mathcal{R}(\eta_Y \circ \mathcal{L}(f)) \circ \varepsilon_X \\
 &= \mathcal{R}(\eta_Y) \circ \mathcal{R} \circ \mathcal{L}(f) \circ \varepsilon_X \\
 &= \mathcal{R}(\eta_Y) \circ \varepsilon_{\mathcal{R}Y} \circ f \\
 &= f
 \end{aligned}$$

and

$$\begin{aligned}
 h^{\mathcal{L}X}(\eta_Y) \circ \mathcal{L}(X, \mathcal{R}Y) \circ h_{\mathcal{R}Y}(\varepsilon_X) \circ \mathcal{R}(\mathcal{L}X, Y)(g) &= \eta_Y \circ \mathcal{L}(\mathcal{R}(g) \circ \varepsilon_X) \\
 &= \eta_Y \circ \mathcal{L} \circ \mathcal{R}(g) \circ \mathcal{L}(\varepsilon_X) \\
 &= g \circ \eta_{\mathcal{L}X} \circ \mathcal{L}(\varepsilon_X) \\
 &= g
 \end{aligned}$$

giving isomorphisms

$$\mathcal{D}(\mathcal{L}X, Y) \cong \mathcal{C}(X, \mathcal{R}Y)$$

for each object X of \mathcal{C} .

Naturality in X follows directly from naturality of ε and η . Namely, given a morphism $f \in \mathcal{C}(X', X)$ we have the commutative diagram

$$\begin{array}{ccc}
 X & \mathcal{D}(\mathcal{L}X, Y) & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{R}Y) \\
 \downarrow f & \downarrow h_{\mathcal{R}Y}(f) & & \downarrow h_Y(\mathcal{L}(f)) \\
 X' & \mathcal{D}(\mathcal{L}X', Y) & \xrightarrow{\sim} & \mathcal{C}(X', \mathcal{R}Y)
 \end{array}$$

$$\begin{array}{ccc}
 \alpha & \xrightarrow{\quad} & \mathcal{R}(\alpha) \circ \varepsilon_X \\
 \downarrow & & \downarrow \\
 \alpha \circ \mathcal{L}(f) & \xrightarrow{\quad} & \mathcal{R}(\alpha \circ \mathcal{L}(f)) \circ \varepsilon_{X'} = \mathcal{R}(\alpha) \circ \varepsilon_X \circ f
 \end{array}$$

and, similarly, the commutative diagram

$$\begin{array}{ccc}
 X & \mathcal{C}(X, \mathcal{R}Y) & \xrightarrow{\sim} & \mathcal{D}(\mathcal{L}X, Y) \\
 \downarrow f & \downarrow h_{\mathcal{R}Y}(f) & & \downarrow h_Y(\mathcal{L}(f)) \\
 X' & \mathcal{C}(X', \mathcal{R}Y) & \xrightarrow{\sim} & \mathcal{D}(\mathcal{L}X', Y)
 \end{array}$$

$$\begin{array}{ccc}
 \beta & \xrightarrow{\quad} & \eta_Y \circ \mathcal{L}(\beta) \\
 \downarrow & & \downarrow \\
 \beta \circ f & \xrightarrow{\quad} & \eta_Y \circ \mathcal{L}(\beta \circ f) = \eta_Y \circ \mathcal{L}(\beta) \circ \mathcal{L}(f)
 \end{array}$$

This establishes the isomorphism of functors.

(iii \Rightarrow i)

First define the functor

$$\mathcal{L}_* : \text{Fun}(\mathcal{D}^{\text{op}}, \mathfrak{Set}) \rightarrow \text{Fun}(\mathcal{C}^{\text{op}}, \mathfrak{Set})$$

as follows. Given $\mathcal{F} : \mathcal{D} \rightarrow \mathfrak{Set}$ and an object X of \mathcal{C} define

$$\mathcal{L}_*(\mathcal{F})(X) = \mathcal{F} \circ \mathcal{L}(X).$$

Given a morphism $\eta \in \text{Fun}(\mathcal{D}^{\text{op}}, \mathfrak{Set})(\mathcal{F}_1, \mathcal{F}_2)$ define natural transformation $\mathcal{L}_*(\eta) : \mathcal{L}_*(\mathcal{F}_1) \rightarrow \mathcal{L}_*(\mathcal{F}_2)$ by

$$\mathcal{L}_*(\eta)_X = \eta_{\mathcal{L}X}.$$

Let $h_-^{\mathcal{D}} : \mathcal{D} \rightarrow \text{Fun}(\mathcal{D}^{\text{op}}, \mathfrak{Set})$ and $h_-^{\mathcal{C}} : \mathcal{C} \rightarrow \text{Fun}(\mathcal{C}^{\text{op}}, \mathfrak{Set})$ denote the Yoneda embeddings, which we note identify \mathcal{C} and \mathcal{D} as full subcategories of $\text{Fun}(\mathcal{D}^{\text{op}}, \mathfrak{Set})$ and $\text{Fun}(\mathcal{C}^{\text{op}}, \mathfrak{Set})$ by the Yoneda Lemma. By assumption, we have for each object Y of \mathcal{D} a representing object $\mathcal{R}Y$ of \mathcal{C} for the functor

$$\mathcal{L}_*(h_Y) = \mathcal{D}(\mathcal{L}(-), Y) \cong \mathcal{C}(-, \mathcal{R}Y) = h_-^{\mathcal{C}}(\mathcal{R}Y)$$

Hence by Lemma 1 we obtain a factorization

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{h_-^{\mathcal{D}}} & \text{Fun}(\mathcal{D}^{\text{op}}, \mathfrak{Set}) & \xrightarrow{\mathcal{L}_*} & \text{Fun}(\mathcal{C}^{\text{op}}, \mathfrak{Set}) \\ & \searrow \exists! \mathcal{R} & & & \nearrow h_-^{\mathcal{C}} \\ & & \mathcal{C} & & \end{array}$$

by the same argument, giving an isomorphism of functors

$$\mathcal{C}(-, \mathcal{R}(-)) = h_-^{\mathcal{C}} \circ \mathcal{R} \cong \mathcal{L}_* \circ h_-^{\mathcal{D}} = \mathcal{D}(\mathcal{L}(-), -),$$

natural in Y .

Similarly, define the functor $\mathcal{R}_* : \text{Fun}(\mathcal{C}, \mathfrak{Set}) \rightarrow \text{Fun}(\mathcal{D}, \mathfrak{Set})$ and note that using the co-Yoneda embeddings $h_{\mathcal{D}}^- : \mathcal{D} \rightarrow \text{Fun}(\mathcal{C}, \mathfrak{Set})$ and $h_{\mathcal{C}}^- : \mathcal{C} \rightarrow \text{Fun}(\mathcal{D}, \mathfrak{Set})$ we obtain a factorization

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{h_{\mathcal{C}}^-} & \text{Fun}(\mathcal{C}, \mathfrak{Set}) & \xrightarrow{\mathcal{R}_*} & \text{Fun}(\mathcal{D}, \mathfrak{Set}) \\ & \searrow \exists! \mathcal{L} & & & \nearrow h_{\mathcal{D}}^- \\ & & \mathcal{D} & & \end{array}$$

giving an isomorphism of functors

$$\mathcal{D}(\mathcal{L}(-), -) = h_{\mathcal{D}}^- \circ \mathcal{L} \cong \mathcal{R}_* \circ h_{\mathcal{C}}^- = \mathcal{C}(-, \mathcal{R}(-)),$$

natural in X . This establishes (i). ■

Proposition 2. *Given an adjunction*

$$\mathcal{C} \begin{array}{c} \xrightarrow{\mathcal{L}} \\ \xleftarrow{\mathcal{R}} \end{array} \mathcal{D}$$

the functor \mathcal{R} commutes with limits and, dually, the functor \mathcal{L} commutes with colimits.

Proof. Let $\mathcal{F}: \mathcal{I} \rightarrow \mathcal{D}$ be a functor for which $\lim \mathcal{F}$ exists. For every morphism $\alpha: i \rightarrow j$ of \mathcal{I} we have a commutative diagram of \mathcal{D}

$$\begin{array}{ccc} & \lim \mathcal{F} & \\ \lambda_i \swarrow & & \searrow \lambda_j \\ \mathcal{F}(i) & \xrightarrow{\mathcal{F}(\alpha)} & \mathcal{F}(j) \end{array}$$

and hence obtain a commutative diagram

$$\begin{array}{ccc} & \mathcal{R}(\lim \mathcal{F}) & \\ \mathcal{R}(\lambda_i) \swarrow & & \searrow \mathcal{R}(\lambda_j) \\ \mathcal{R} \circ \mathcal{F}(i) & \xrightarrow{\mathcal{R} \circ \mathcal{F}(\alpha)} & \mathcal{R} \circ \mathcal{F}(j) \end{array}$$

of \mathcal{C} .

Given any other object Z of \mathcal{C} equipped with morphisms $\zeta_i: Z \rightarrow \mathcal{F}(i)$ such that $\mathcal{R} \circ \mathcal{F}(\alpha) \circ \zeta_i = \zeta_j$ for all $\alpha \in \mathcal{I}(i, j)$ for all i, j , we obtain by the natural isomorphism

$$\begin{array}{ccc} \mathcal{C}(Z, \mathcal{R} \circ \mathcal{F}(i)) & \xrightarrow{\mathcal{L}(Z, \mathcal{R} \circ \mathcal{F}(i))} & \mathcal{D}(\mathcal{L}Z, \mathcal{L} \circ \mathcal{R} \circ \mathcal{F}(i)) & \xrightarrow{h^{\mathcal{L}Z}(\varepsilon_{\mathcal{F}(i)})} & \mathcal{D}(\mathcal{L}Z, \mathcal{F}(i)) \\ \zeta_i \dashv \vdash & \longrightarrow & \mathcal{L}(\zeta_i) \dashv \vdash & \longrightarrow & \mathcal{L}(\zeta_i) \circ \varepsilon_{\mathcal{F}(i)} \end{array}$$

for all objects i of \mathcal{I} . By the naturality diagram

$$\begin{array}{ccc} \mathcal{C}(Z, \mathcal{R} \circ \mathcal{F}(i)) & \xrightarrow{\sim} & \mathcal{D}(\mathcal{R}Z, \mathcal{F}(i)) \\ \downarrow h^Z(\mathcal{R} \circ \mathcal{F}(\alpha)) & & \downarrow h^{\mathcal{R}Z}(\mathcal{F}(\alpha)) \\ \mathcal{C}(Z, \mathcal{R} \circ \mathcal{F}(j)) & \xrightarrow{\sim} & \mathcal{D}(\mathcal{R}Z, \mathcal{F}(j)) \end{array}$$

we see that for all $\alpha \in \mathcal{I}(i, j)$ over all objects i, j of \mathcal{I}

$$\mathcal{F}(\alpha) \circ \mathcal{L}(\zeta_i) \circ \varepsilon_{\mathcal{F}(i)} = \mathcal{L}(\mathcal{R} \circ \mathcal{F}(\alpha) \circ \zeta_i) \circ \varepsilon_{\mathcal{F}(j)} = \mathcal{L}(\zeta_j) \circ \varepsilon_{\mathcal{F}(j)}$$

and hence we obtain a unique morphism $h \in \mathcal{D}(\mathcal{R}Z, \lim \mathcal{F})$ making all diagrams

$$\begin{array}{ccc} & \mathcal{L}Z & \\ \mathcal{L}(\zeta_i) \circ \varepsilon_{\mathcal{F}(i)} \swarrow & \downarrow \exists! h & \searrow \mathcal{L}(\zeta_j) \circ \varepsilon_{\mathcal{F}(j)} \\ & \lim \mathcal{F} & \\ \lambda_i \swarrow & & \searrow \lambda_j \\ \mathcal{F}(i) & \xrightarrow{\mathcal{F}(\alpha)} & \mathcal{F}(j) \end{array}$$

indexed over $\alpha \in \mathcal{I}(i, j)$ commute. Pulling this map back along the isomorphism

$$\begin{array}{ccc} \mathcal{D}(\mathcal{L}Z, \lim \mathcal{F}) & \longrightarrow & \mathcal{C}(\mathcal{R} \circ \mathcal{L}Z, \mathcal{R} \lim \mathcal{F}) & \longrightarrow & \mathcal{C}(Z, \mathcal{R} \lim \mathcal{F}) \\ h \dashv \vdash & \longrightarrow & \mathcal{R}(h) \dashv \vdash & \longrightarrow & \eta_Z \circ \mathcal{R}(h) \end{array}$$

gives a morphism, which we see makes all diagrams

$$\begin{array}{ccc}
 & Z & \\
 \zeta_i \swarrow & \downarrow \eta_Z \circ \mathcal{R}(h) & \searrow \zeta_j \\
 \mathcal{R} \lim \mathcal{F} & & \\
 \mathcal{R}(\lambda_i) \swarrow & & \searrow \mathcal{R}(\lambda_j) \\
 \mathcal{R} \circ \mathcal{F}(i) & \xrightarrow{\mathcal{R} \circ \mathcal{F}(\alpha)} & \mathcal{R} \circ \mathcal{F}(j)
 \end{array}$$

commute by chasing h through the triangular prism

$$\begin{array}{ccccc}
 & \mathcal{D}(\mathcal{L}Z, \lim \mathcal{F}) & & & \\
 & \downarrow \wr & & & \\
 & \mathcal{C}(Z, \mathcal{R} \lim \mathcal{F}) & & & \\
 \begin{array}{c} \swarrow \\ \downarrow \\ \searrow \end{array} & & & & \begin{array}{c} \swarrow \\ \downarrow \\ \searrow \end{array} \\
 \mathcal{D}(\mathcal{L}Z, \mathcal{F}(i)) & \xrightarrow{h^{\mathcal{L}Z}(\mathcal{F}(\alpha))} & \mathcal{D}(\mathcal{L}Z, \mathcal{F}(j)) & & \\
 \downarrow \wr & \begin{array}{c} \swarrow \\ \downarrow \\ \searrow \end{array} & & & \downarrow \wr \\
 \mathcal{C}(Z, \mathcal{R} \circ \mathcal{F}(i)) & \xrightarrow{h^Z(\mathcal{R} \circ \mathcal{F}(\alpha))} & \mathcal{C}(Z, \mathcal{R} \circ \mathcal{F}(j)) & &
 \end{array}$$

which commutes since the left, right, and bottom faces are naturality squares and the rear face commutes by construction of h .

We note that $\eta_Z \circ \mathcal{R}(h)$ is necessarily unique, for if we had some other morphism $h' \in \mathcal{C}(Z, \mathcal{R} \lim \mathcal{F})$ making the relevant diagrams of \mathcal{C} commute, then pushing it along the isomorphism gives $\mathcal{L}(h') \circ \varepsilon_{\lim \mathcal{F}} \in \mathcal{D}(\mathcal{L}Z, \lim \mathcal{F})$ making the relevant diagrams of \mathcal{D} commute, which implies $\mathcal{L}(h') \circ \varepsilon_{\lim \mathcal{F}} = h$ by unicity, and hence $h' = \eta_Z \circ \mathcal{R}(h)$. Therefore $\mathcal{R} \lim \mathcal{F} \cong \lim \mathcal{R} \circ \mathcal{F}$ by unicity of universals. \blacksquare

Proposition 3. *Let $\mathcal{C}, \mathcal{D}, \mathcal{D}'$ be categories equipped with adjunctions*

$$\mathcal{C} \begin{array}{c} \xrightarrow{\mathcal{L}} \\ \xleftarrow{\mathcal{R}} \end{array} \mathcal{D} \quad \text{and} \quad \mathcal{D} \begin{array}{c} \xrightarrow{\mathcal{L}'} \\ \xleftarrow{\mathcal{R}'} \end{array} \mathcal{D}'$$

Then

$$\mathcal{C} \begin{array}{c} \xrightarrow{\mathcal{L}' \circ \mathcal{L}} \\ \xleftarrow{\mathcal{R} \circ \mathcal{R}'} \end{array} \mathcal{D}'$$

is an adjunction.

Proof. The composition of functors is clearly well defined, so we obtain a natural isomorphism

$$\begin{aligned}
 \mathcal{D}'(\mathcal{L}' \circ \mathcal{L}(X), Y') &\cong \mathcal{D}(\mathcal{L}(X), \mathcal{R}'Y') \\
 &\cong \mathcal{C}(X, \mathcal{R} \circ \mathcal{R}'Y').
 \end{aligned}$$

\blacksquare

Definition 1. Given two categories \mathcal{C} and \mathcal{D} , an *equivalence of categories* is a pair of functors $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{G}: \mathcal{D} \rightarrow \mathcal{C}$ and natural isomorphisms $\varepsilon: \text{id}_{\mathcal{C}} \rightarrow \mathcal{G} \circ \mathcal{F}$, $\eta: \mathcal{F} \circ \mathcal{G} \rightarrow \text{id}_{\mathcal{D}}$.

Proposition 4. Let $\mathcal{L}: \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{R}: \mathcal{C} \rightarrow \mathcal{D}$ be adjoint functors with unit $\varepsilon: \text{id}_{\mathcal{C}} \rightarrow \mathcal{R} \circ \mathcal{L}$ and counit $\eta: \mathcal{L} \circ \mathcal{R} \rightarrow \text{id}_{\mathcal{D}}$.

- (i) The functor \mathcal{R} is fully faithful if and only if η is an isomorphism,
- (ii) The functor \mathcal{L} is fully faithful if and only if ε is an isomorphism
- (iii) The following are equivalent:
 - (a) \mathcal{L} is an equivalence of categories,
 - (b) \mathcal{R} is an equivalence of categories,
 - (c) \mathcal{L} and \mathcal{R} are fully faithful. In this case, \mathcal{L} and \mathcal{R} are quasi-inverses of one another, and ε , η are both isomorphism.

Proof. For (i), we observe from the diagram

$$\begin{array}{ccc} \mathcal{D}(Y, Y') & \xrightarrow{\mathcal{R}(Y, Y')} & \mathcal{C}(\mathcal{R}Y, \mathcal{R}Y') \\ & \searrow h_{Y'}(\eta_Y) & \downarrow \wr \\ & & \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y') \end{array}$$

that $\mathcal{R}(Y, Y')$ is an isomorphism if and only if $h_{Y'}(\eta_Y)$ is an isomorphism. Hence \mathcal{R} is fully faithful if and only if $h_{Y'}(\eta_Y)$ is an isomorphism for all objects Y, Y' of \mathcal{D} , and we are reduced to showing this is equivalent to η_Y being an isomorphism.

Clearly if η_Y is an isomorphism, then $h_{Y'}(\eta_Y)$ is an isomorphism for all objects Y, Y' of \mathcal{D} with inverse

$$h_{Y'}(\eta_Y^{-1}): \mathcal{D}(\mathcal{L} \circ \mathcal{R}Y, Y').$$

Conversely, taking $Y' = Y$ we obtain an inverse to η_Y by applying the the inverse morphism:

$$\eta_Y^{-1} = h_Y(\eta_Y)^{-1}(\eta_Y).$$

For (ii), we apply the same argument, *mutatis mutandis*, to the diagram

$$\begin{array}{ccc} \mathcal{C}(X, X') & \xrightarrow{\mathcal{L}(X, X')} & \mathcal{D}(\mathcal{L}X, \mathcal{L}X') \\ & \searrow h^X(\varepsilon_X) & \downarrow \wr \\ & & \mathcal{C}(X, \mathcal{R} \circ \mathcal{L}X') \end{array}$$

Part (iii) follows from the definition of an equivalence of categories and the unicity of the adjoints. ■

3. QUILLEN ADJUNCTIONS AND DERIVED FUNCTORS

Definition 2. Let \mathcal{C} and \mathcal{D} be model categories.

- (1) A functor $\mathcal{L}: \mathcal{C} \rightarrow \mathcal{D}$ is *left Quillen* if \mathcal{L} is a left adjoint and preserves cofibrations and trivial cofibrations.
- (2) A functor $\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$ is *right Quillen* if \mathcal{R} is a right adjoint and preserves fibrations and trivial fibrations.

(3) An adjunction $(\mathcal{L}, \mathcal{R}, \varphi)$, where φ is the natural isomorphism

$$\varphi_{X,Y}: \mathcal{D}(\mathcal{L}X, Y) \xrightarrow{\sim} \mathcal{C}(X, \mathcal{R}Y),$$

is called a *Quillen adjunction* if \mathcal{L} is left Quillen.

Remark 1. Hovey always uses η for the unit of adjunction and ε for the counit of adjunction; my notation is exactly the opposite.

Example 1. Let \mathcal{C} be a model category and let I be a set. Equip the product category \mathcal{C}^I with the product model structure. We can view an object of \mathcal{C}^I as a discrete diagram of \mathcal{C} and a morphism of this category as a morphism of discrete diagrams. For example, in the case where I has two elements, a morphism between two objects, $f = (f_1, f_2): X = (X_1, X_2) \rightarrow Y = (Y_1, Y_2)$ is just two morphisms of \mathcal{C} :

$$\begin{array}{ccc} X_1 & & X_2 \\ \downarrow f_1 & & \downarrow f_2 \\ Y_1 & & Y_2 \end{array}$$

with no commutativity relations.

As \mathcal{C} has all small limits by assumption, we can define a product functor $\lim_I: \mathcal{C}^I \rightarrow \mathcal{C}$ which takes an object $X = (X_i)_I$ of \mathcal{C}^I to the limit over the discrete diagram,

$$\lim_I X = \lim_I X_i = \prod_{i \in I} X_i.$$

For any morphism of $f \in \mathcal{C}^I(X, Y)$, $\lim_I(f)$ is the unique map induced by the universal property of limit, which is determined by the diagrams

$$\begin{array}{ccc} \lim_I X & \xrightarrow{\lim_I(f)} & \lim_I Y \\ \downarrow & & \downarrow \\ X_i & \xrightarrow{f_i} & Y_i \end{array}$$

indexed over I . For example, when I has two elements, the image of a morphism $(f_1, f_2): (X_1, X_2) \rightarrow (Y_1, Y_2)$ is just the product map,

$$\begin{array}{ccccc} X_1 & \longleftarrow & X_1 \times X_2 & \longrightarrow & X_2 \\ \downarrow f_1 & & \downarrow \exists! f_1 \times f_2 & & \downarrow f_2 \\ Y_1 & \longleftarrow & Y_1 \times Y_2 & \longrightarrow & Y_2 \end{array}$$

We can define a diagonal functor $\Delta: \mathcal{C} \rightarrow \mathcal{C}^I$ which takes an object X of \mathcal{C} to the discrete diagram with I copies of X . A morphism $f \in \mathcal{C}(X, X)$ defines a morphism of discrete diagrams by taking f in each component. We note that, by definition, a morphism $f \in \mathcal{C}^I(\Delta X, Y)$ is simply a collection of morphisms $f_i \in \mathcal{C}(X, Y_i)$, giving

$$\mathcal{C}^I(\Delta X, Y) = \prod_{i \in I} \mathcal{C}(X, Y_i)$$

With this observation, it's easy to see that this functor is left adjoint to \lim_I since we can view the universal property of products as a representability statement: The object $\lim_I Y_i = \prod_{i \in I} Y_i$ is the representing object

of the functor $\prod_{i \in I} \mathcal{C}(-, Y_i)$, and hence

$$\mathcal{C}\left(-, \lim_I Y\right) = \mathcal{C}\left(-, \prod_{i \in I} Y_i\right) \cong \prod_{i \in I} \mathcal{C}(-, Y_i) = \mathcal{C}^I(-, Y) = \mathcal{C}^I(\Delta(-), Y).$$

By Proposition 1 we obtain the desired natural isomorphism

$$\mathcal{C}^I(\Delta X, Y) = \prod_{i \in I} \mathcal{C}(\Delta X, Y_i) \cong \mathcal{C}\left(\Delta X, \lim_I Y\right).$$

Note that by definition of the product model structure, Δ necessarily preserves cofibrations and weak equivalences, so this adjunction is Quillen.

Lemma 2. *Let \mathcal{C} and \mathcal{D} be model categories and suppose $\mathcal{L}: \mathcal{C} \rightarrow \mathcal{D}$ is left adjoint to $\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$. This adjunction is Quillen if and only if \mathcal{R} is right Quillen.*

Proof. By Lemma 1.1.10 of Hovey it suffices to show that given a fibration (resp. trivial fibration) $p \in \mathcal{D}(Y_1, Y_2)$, then $\mathcal{R}p \in \mathcal{C}(\mathcal{R}Y_1, \mathcal{R}Y_2)$ has the right lifting property with respect to all trivial cofibrations (resp. cofibrations). Let $f \in \mathcal{C}(X_1, X_2)$ be a trivial cofibration (resp. cofibration). To say that $\mathcal{R}p$ has the right lifting property with respect to f is to say that for every commutative diagram

$$\begin{array}{ccc} X_1 & \xrightarrow{\varphi} & \mathcal{R}Y_1 \\ \downarrow f & & \downarrow \mathcal{R}p \\ X_2 & \xrightarrow{\psi} & \mathcal{R}Y_2 \end{array}$$

there is a lift $\ell: X_2 \rightarrow \mathcal{R}Y_1$ such that $\ell \circ f = \varphi$ and $\mathcal{R}p \circ \ell = \psi$. This is equivalent to the morphisms of sets

$$\mathcal{D}(\mathcal{L}X_2, Y_1) \cong \mathcal{C}(X_2, \mathcal{R}Y_1) \xrightarrow{h^{X_2}(\mathcal{R}p)} \mathcal{C}(X_2, \mathcal{R}Y_2) \cong \mathcal{D}(\mathcal{L}X_2, Y_2)$$

and

$$\mathcal{D}(\mathcal{L}X_2, Y_1) \cong \mathcal{C}(X_2, \mathcal{R}Y_1) \xrightarrow{h^{\mathcal{R}Y_1}(f)} \mathcal{C}(X_1, \mathcal{R}Y_1) \cong \mathcal{D}(\mathcal{L}X_1, Y_1)$$

being surjective. However, by the naturality squares

$$\begin{array}{ccc} \mathcal{D}(\mathcal{L}X_2, Y_1) & \xrightarrow{\sim} & \mathcal{C}(X_2, \mathcal{R}Y_1) \\ \downarrow h^{\mathcal{L}X_2}(p) & & \downarrow h^{X_2}(\mathcal{R}p) \\ \mathcal{D}(\mathcal{L}X_2, Y_2) & \xrightarrow{\sim} & \mathcal{C}(X_2, \mathcal{R}Y_2) \end{array} \quad \text{and} \quad \begin{array}{ccc} \mathcal{D}(\mathcal{L}X_2, Y_1) & \xrightarrow{\sim} & \mathcal{C}(X_2, \mathcal{R}Y_1) \\ \downarrow h_{Y_1}(\mathcal{L}(f)) & & \downarrow h^{\mathcal{R}Y_1}(f) \\ \mathcal{D}(\mathcal{L}X_1, Y_1) & \xrightarrow{\sim} & \mathcal{C}(X_1, \mathcal{R}Y_1) \end{array}$$

this is equivalent to $\mathcal{L}(f)$ having the left lifting property with respect to p . This is equivalent to $\mathcal{L}(f)$ being a cofibration (resp. trivial cofibration), which in turn is equivalent to the adjunction being Quillen. \blacksquare

Remark 2. (i) If $\mathcal{L}: \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$ is a Quillen adjunction, then, by Ken Brown's Lemma, \mathcal{L} preserves weak equivalences between cofibrant objects and \mathcal{R} preserves weak equivalences between fibrant objects. By abuse of notation, we write $\mathcal{L}: \mathcal{C}_c \rightarrow \mathcal{D}$ and $\mathcal{R}: \mathcal{D}_f \rightarrow \mathcal{C}$ for the restrictions, then observe that we obtain

$$\begin{array}{ccccc} \mathcal{C} & \xrightarrow{Q} & \mathcal{C}_c & \xrightarrow{\mathcal{L}} & \mathcal{D} \\ \downarrow \gamma_{\mathcal{C}} & & \downarrow \gamma_{\mathcal{C}_c} & \exists! \nearrow & \downarrow \gamma_{\mathcal{D}} \\ \text{Ho}(\mathcal{C}) & \xrightarrow{\text{Ho}(Q)} & \text{Ho}(\mathcal{C}_c) & \xrightarrow{\text{Ho}(\mathcal{L})} & \text{Ho}(\mathcal{D}) \end{array}$$

and

$$\begin{array}{ccccc}
\mathcal{D} & \xrightarrow{R} & \mathcal{D}_f & \xrightarrow{\mathcal{R}} & \mathcal{C} \\
\downarrow \gamma_{\mathcal{D}} & & \downarrow \gamma_{\mathcal{D}_f} & \dashrightarrow \exists! & \downarrow \gamma_{\mathcal{C}} \\
\mathrm{Ho}(\mathcal{D}) & \xrightarrow{\mathrm{Ho}(R)} & \mathrm{Ho}(\mathcal{D}_f) & \xrightarrow{\mathrm{Ho}(\mathcal{R})} & \mathrm{Ho}(\mathcal{C})
\end{array}$$

by the universal property for the homotopy category.

We also note that given natural transformations $\eta: \mathcal{L} \rightarrow \mathcal{L}'$, $\nu: \mathcal{R} \rightarrow \mathcal{R}'$ between Quillen functors, we obtain natural transformations $\mathrm{Ho}(\eta): \mathrm{Ho}(\mathcal{L}) \rightarrow \mathrm{Ho}(\mathcal{L}')$ and $\mathrm{Ho}(\nu): \mathrm{Ho}(\mathcal{R}) \rightarrow \mathrm{Ho}(\mathcal{R}')$ defined by $\mathrm{Ho}(\eta)_X = \eta_X$ and $\mathrm{Ho}(\nu)_X = \nu_X$. Since all the functors involved preserve weak equivalences, these are indeed natural by Lemma 1.2.2.

Note that on objects,

$$\mathrm{Ho}(\mathcal{L})(X) = \mathcal{L}(QX) \text{ and } \mathrm{Ho}(\mathcal{R})(Y) = \mathcal{R}(QY).$$

- (ii) This construction does not require Quillen functors. We only used the fact that these functors preserve weak equivalences between cofibrant (resp. fibrant) objects.

Definition 3. Let \mathcal{C} and \mathcal{D} be model categories.

- (1) If $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$ is a left Quillen functor, define the *total left derived functor* $L\mathcal{F}: \mathrm{Ho}(\mathcal{C}) \rightarrow \mathrm{Ho}(\mathcal{D})$ to be the composition

$$\mathrm{Ho}(\mathcal{C}) \xrightarrow{\mathrm{Ho}(Q)} \mathrm{Ho}(\mathcal{C}_c) \xrightarrow{\mathrm{Ho}(\mathcal{F})} \mathrm{Ho}(\mathcal{D})$$

Given a natural transformation $\tau: \mathcal{F} \rightarrow \mathcal{F}'$ of left Quillen functors, define the *total derived natural transformation* $L\tau$ to be $\mathrm{Ho}(\tau) \circ \mathrm{Ho}(Q)$, so that $(L\tau)_X = \tau_{QX}$.

- (2) If $\mathcal{G}: \mathcal{D} \rightarrow \mathcal{C}$ is a right Quillen functor, define the *total right derived functor* $R\mathcal{G}: \mathrm{Ho}(\mathcal{D}) \rightarrow \mathrm{Ho}(\mathcal{C})$ to be the composition

$$\mathrm{Ho}(\mathcal{D}) \xrightarrow{\mathrm{Ho}(R)} \mathrm{Ho}(\mathcal{D}_f) \xrightarrow{\mathrm{Ho}(\mathcal{G})} \mathrm{Ho}(\mathcal{C})$$

Given a natural transformation $\tau: \mathcal{G} \rightarrow \mathcal{G}'$ of right Quillen functors, define the *total derived natural transformation* $R\tau$ to be $\mathrm{Ho}(\tau) \circ \mathrm{Ho}(R)$, so that $(R\tau)_X = \tau_{RX}$.

Theorem 2. For every model category, \mathcal{C} , there is a natural isomorphism $\alpha: L(\mathrm{id}_{\mathcal{C}}) \rightarrow \mathrm{id}_{\mathrm{Ho}(\mathcal{C})}$. Also for every pair of left Quillen functors $F: \mathcal{C} \rightarrow \mathcal{D}$ and $F': \mathcal{D} \rightarrow \mathcal{E}$, there is a natural isomorphism $m = m_{\mathcal{F}', \mathcal{F}}: L\mathcal{F}' \circ L\mathcal{F} \rightarrow L(\mathcal{F}' \circ \mathcal{F})$. These natural isomorphisms satisfy the following properties.

- (1) An associativity coherence diagram is commutative. That is, if $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{C}'$, $\mathcal{F}': \mathcal{C}' \rightarrow \mathcal{C}''$, and $\mathcal{F}'': \mathcal{C}'' \rightarrow \mathcal{C}'''$ are left Quillen functors, then the diagram

$$\begin{array}{ccccc}
(L\mathcal{F}'' \circ L\mathcal{F}') \circ L\mathcal{F} & \xrightarrow{m_{\mathcal{F}'', \mathcal{F}' \circ L\mathcal{F}}} & L(\mathcal{F}'' \circ \mathcal{F}') \circ L\mathcal{F} & \xrightarrow{m_{(\mathcal{F}'' \circ \mathcal{F}'), \mathcal{F}}} & L((\mathcal{F}'' \circ \mathcal{F}') \circ \mathcal{F}) \\
\parallel & & & & \parallel \\
L\mathcal{F}'' \circ (L\mathcal{F}' \circ L\mathcal{F}) & \xrightarrow{L\mathcal{F}'' \circ m_{\mathcal{F}', \mathcal{F}}} & L\mathcal{F}'' \circ L(\mathcal{F}' \circ \mathcal{F}) & \xrightarrow{m_{\mathcal{F}'', (\mathcal{F}' \circ \mathcal{F})}} & L(\mathcal{F}'' \circ (\mathcal{F}' \circ \mathcal{F}))
\end{array}$$

commutes.

(2) A left unit coherence diagram is commutative. That is, if $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$ is a left Quillen functor, then the diagram

$$\begin{array}{ccc} L(\text{id}_{\mathcal{D}}) \circ L\mathcal{F} & \xrightarrow{m} & L(\text{id}_{\mathcal{D}} \circ \mathcal{F}) \\ \downarrow \alpha \circ L\mathcal{F} & & \parallel \\ 1_{\text{Ho}(\mathcal{D})} \circ L\mathcal{F} & \xlongequal{\quad} & L\mathcal{F} \end{array}$$

commutes.

(3) A right unit coherence diagram is commutative. That is, if $\mathcal{F}: \mathcal{C} \rightarrow \mathcal{D}$ is a left Quillen functor, then the diagram

$$\begin{array}{ccc} L\mathcal{F} \circ L(\text{id}_{\mathcal{C}}) & \xrightarrow{m} & L(\mathcal{F} \circ \text{id}_{\mathcal{C}}) \\ \downarrow L\mathcal{F} \circ \alpha & & \parallel \\ L\mathcal{F} \circ \text{id}_{\text{Ho}(\mathcal{C})} & \xlongequal{\quad} & L\mathcal{F} \end{array}$$

commutes.

Proof. Let $\iota: \mathcal{C}_c \rightarrow \mathcal{C}$ be the natural inclusion. We observe from the functorial factorization

$$\begin{array}{ccc} 0 & \xrightarrow{\quad} & X \\ & \searrow & \nearrow q_X \\ & & QX \end{array}$$

that we obtain a natural transformation $q: \iota \circ Q \rightarrow \text{id}_{\mathcal{C}}$. Taking the derived natural transformation $\text{Ho}(q)$ gives the natural isomorphism

$$\text{id}_{\text{Ho}(\mathcal{C})} \xrightarrow{\sim} \text{Ho}(\iota) \circ \text{Ho}(Q) \xrightarrow{\text{Ho}(q)} L(\text{id}_{\mathcal{C}})$$

because q_X is a trivial fibration.

Define $m_{\mathcal{F}', \mathcal{F}}$ to be the collection of maps

$$L\mathcal{F}' \circ L\mathcal{F}X = \mathcal{F}'(Q(\mathcal{F}(QX))) \xrightarrow{F'(q_{\mathcal{F}(QX)})} \mathcal{F}'(\mathcal{F}(QX)) = L(\mathcal{F}' \circ \mathcal{F}(X))$$

which is natural in X as a functor on \mathcal{C} , for if $f' \in \mathcal{C}(X, X')$ then we have the commutative diagram

$$\begin{array}{ccc} QX & \xrightarrow{q_X} & X \\ \downarrow Qf & & \downarrow f \\ QX' & \xrightarrow{q_{X'}} & X' \end{array}$$

which gives rise to the commutative diagram in \mathcal{D}

$$\begin{array}{ccccc} Q(\mathcal{F}(QX)) & \xrightarrow{q_{\mathcal{F}(QX)}} & \mathcal{F}(QX) & \xrightarrow{\mathcal{F}(q_X)} & \mathcal{F}(X) \\ \downarrow Q\mathcal{F}(Qf) & & \downarrow \mathcal{F}(Qf) & & \downarrow \mathcal{F}(f) \\ Q(\mathcal{F}(QX')) & \xrightarrow{q_{\mathcal{F}(QX')}} & \mathcal{F}(QX') & \xrightarrow{\mathcal{F}(q_{X'})} & \mathcal{F}(X') \end{array}$$

and hence a commutative diagram in \mathcal{E}

$$\begin{array}{ccccc}
\mathcal{F}'(Q(\mathcal{F}(QX))) & \xrightarrow{\mathcal{F}'(q_{\mathcal{F}(QX)})} & \mathcal{F}' \circ \mathcal{F}(QX) & \xrightarrow{\mathcal{F}' \circ \mathcal{F}(q_X)} & \mathcal{F}' \circ \mathcal{F}(X) \\
\downarrow \mathcal{F}'(Q\mathcal{F}(Qf)) & & \downarrow \mathcal{F}' \circ \mathcal{F}(Qf) & & \downarrow \mathcal{F}' \circ \mathcal{F}(f) \\
\mathcal{F}'(Q(\mathcal{F}(QX'))) & \xrightarrow{\mathcal{F}'(q_{\mathcal{F}(QX')})} & \mathcal{F}' \circ \mathcal{F}(QX') & \xrightarrow{\mathcal{F}' \circ \mathcal{F}(q_{X'})} & \mathcal{F}' \circ \mathcal{F}(X')
\end{array}$$

Since all functors involved preserve weak equivalences, $m_{\mathcal{F}', \mathcal{F}}$ is also natural in X as a functor on $\text{Ho}(\mathcal{C})$. Moreover, \mathcal{F} preserves cofibrant objects because it preserves cofibrations as a left Quillen functor, hence \mathcal{F}' preserves the weak equivalence $q_{\mathcal{F}(QX)}$ between cofibrant objects and thus it follows that $m_{\mathcal{F}', \mathcal{F}}$ is an isomorphism in $\text{Ho}(\mathcal{E})$.

For the associativity coherence diagram, we must show that

$$(\mathcal{F}'' \circ \mathcal{F}'(q_{\mathcal{F}QX})) \circ (F''(q_{\mathcal{F}', Q\mathcal{F}QX})) = (F''(q_{\mathcal{F}', \mathcal{F}QX})) \circ ((\mathcal{F}'' \circ Q \circ \mathcal{F}'(q_{\mathcal{F}QX})).$$

This follows from naturality of q and a construction similar to the one above.

The left unit coherence diagram commutes because by definition

$$m_{\text{id}_{\mathcal{D}}, \mathcal{F}}(X) = \text{id}_{\mathcal{D}}(q_{\mathcal{F}QX}) = q_{\mathcal{F}QX}$$

and

$$\alpha \circ L\mathcal{F}(X) = \alpha_{\mathcal{F}QX} = q_{\mathcal{F}QX}$$

For the right unit coherence diagram, we have

$$m_{\mathcal{F}, \text{id}_{\mathcal{E}}}(X) = \mathcal{F}(q_{\text{id}_{\mathcal{E}}(q_{QX})}) = \mathcal{F}(q_{QX}): \mathcal{F}(QQX) \rightarrow \mathcal{F}(QX)$$

and

$$L\mathcal{F} \circ \alpha(X) = \mathcal{F}(Q(q_X)): \mathcal{F}(QQX) \rightarrow \mathcal{F}(QX).$$

Given a cofibrant object X we have the commutative diagram

$$\begin{array}{ccc}
QQX & \xrightarrow{q_{QX}} & QX \\
\downarrow Q(q_X) & & \downarrow q_X \\
QX & \xrightarrow{q_X} & X
\end{array}$$

because q is natural. Since q_X is a weak equivalence between cofibrant objects it follows that $\mathcal{F}(q_X)$ is invertible in $\text{Ho}(\mathcal{D})$ and hence

$$\mathcal{F}(q_{QX}) = \mathcal{F}Q(q_X).$$

Since every object of $\text{Ho}(\mathcal{C})$ is weakly equivalent to a cofibrant object, this completes the proof. \blacksquare

Definition 4. Let \mathcal{C} , \mathcal{D} , and \mathcal{E} be categories and let $\mathcal{F}, \mathcal{G}: \mathcal{C} \rightarrow \mathcal{D}$, $\mathcal{F}', \mathcal{G}': \mathcal{D} \rightarrow \mathcal{E}$ be functors. Given natural transformations $\eta: \mathcal{F} \rightarrow \mathcal{G}$ and $\nu: \mathcal{F}' \rightarrow \mathcal{G}'$, define the *horizontal composition* $\eta * \nu: \mathcal{F}' \circ \mathcal{F} \rightarrow \mathcal{G}' \circ \mathcal{G}$ is the natural transformation defined by the collection of morphisms of \mathcal{E}

$$\begin{array}{ccc}
\mathcal{F}' \circ \mathcal{F}(X) & \xrightarrow{\mathcal{F}'(\eta_X)} & \mathcal{F}' \circ \mathcal{G}(X) \\
\downarrow \nu_{\mathcal{F}X} & & \downarrow \nu_{\mathcal{G}X} \\
\mathcal{G}' \circ \mathcal{F}(X) & \xrightarrow{\mathcal{G}'(\eta_X)} & \mathcal{G}' \circ \mathcal{G}(X)
\end{array}$$

Lemma 3. Let \mathcal{C} , \mathcal{D} , and \mathcal{E} be model categories. Let $\mathcal{F}, \mathcal{G}: \mathcal{C} \rightarrow \mathcal{D}$ and $\mathcal{F}', \mathcal{G}': \mathcal{D} \rightarrow \mathcal{E}$ be left Quillen functors. Suppose $\eta: \mathcal{F} \rightarrow \mathcal{G}$ and $\nu: \mathcal{F}' \rightarrow \mathcal{G}'$ be natural transformations. If m is the composition isomorphism of the Theorem above, then the diagram

$$\begin{array}{ccc} L\mathcal{F}' \circ L\mathcal{F} & \xrightarrow{m} & L(\mathcal{F}' \circ \mathcal{F}) \\ \downarrow L\eta * L\nu & & \downarrow L(\eta * \nu) \\ L\mathcal{G}' \circ L\mathcal{G} & \xrightarrow{m} & L(\mathcal{G}' \circ \mathcal{G}) \end{array}$$

commutes.

Proof. We unravel the definitions. The morphism

$$L(\eta * \nu) \circ m_X : \mathcal{F}' Q \mathcal{F} Q X \rightarrow \mathcal{G}' \mathcal{G} Q X$$

is given by the composition across the top of the commutative diagram

$$\begin{array}{ccccc} \mathcal{F}' Q \mathcal{F} Q X & \xrightarrow{\mathcal{F}'(q_{\mathcal{F} Q X})} & \mathcal{F}' \mathcal{F} Q X & \xrightarrow{\nu_{\mathcal{F} Q X}} & \mathcal{G}' \mathcal{F} Q X \\ \downarrow \mathcal{F}' Q(\eta_{Q X}) & & \downarrow \mathcal{F}'(\eta_{Q X}) & & \downarrow \mathcal{G}'(\nu_{Q X}) \\ \mathcal{F}' Q \mathcal{G} Q X & \xrightarrow{\mathcal{F}'(q_{\mathcal{G} Q X})} & \mathcal{F}' \mathcal{G} Q X & \xrightarrow{\nu_{\mathcal{G} Q X}} & \mathcal{G}' \mathcal{G} Q X \\ \downarrow \nu_{Q \mathcal{G} Q X} & & \downarrow \nu_{\mathcal{G} Q X} & & \parallel \\ \mathcal{G}' Q \mathcal{G} Q X & \xrightarrow{\mathcal{G}'(q_{\mathcal{G} Q X})} & \mathcal{G}' \mathcal{G} Q X & \xlongequal{\quad} & \mathcal{G}' \mathcal{G} Q X \end{array}$$

and the morphism $m \circ L\eta * L\nu_X : \mathcal{F}' Q \mathcal{F} Q X \rightarrow \mathcal{G}' \mathcal{G} Q X$ is given by the composition across the bottom of the commutative diagram

$$\begin{array}{ccccc} \mathcal{F}' Q \mathcal{F}(Q X) & \xrightarrow{\nu_{Q \mathcal{F} Q X}} & \mathcal{G}' Q \mathcal{F} Q X & \xrightarrow{\mathcal{G}'(q_{\mathcal{F} Q X})} & \mathcal{G}' \mathcal{F} Q X \\ \downarrow \mathcal{F}' Q(\eta_{Q X}) & & \downarrow \mathcal{G}' Q(\eta_{Q X}) & & \downarrow \mathcal{G}'(\eta_{Q X}) \\ \mathcal{F}' Q \mathcal{G} Q X & \xrightarrow{\nu_{Q \mathcal{G} Q X}} & \mathcal{G}' Q \mathcal{G} Q X & \xrightarrow{\mathcal{G}'(q_{\mathcal{G} Q X})} & \mathcal{G}' \mathcal{G} Q X \end{array}$$

Chasing the bottom left side of each diagram gives us the desired equality. \blacksquare

Remark 3. Essentially, this says we have a 2-category (modulo some set theoretic issues...) with 0-cells model categories, 1-cells left Quillen functors, 2-cells the natural transformations and the homotopy category, total derived functor, and total derived natural transformation define a pseudo 2-functor to the 2-category of categories.

Lemma 4. Let \mathcal{C} be a model category and assume we have homotopic morphisms $f_0 \sim f_1 \in \mathcal{C}(X, Y)$. If Y is fibrant object, then we may always choose a fibrant path object. Dually, if X cofibrant we may always choose a cofibrant cylinder object.

Proof. Since Y is assumed to be fibrant, the pullback

$$\begin{array}{ccc} Y \times Y & \xrightarrow{\pi_0} & Y \\ \downarrow \pi_1 & & \downarrow \\ Y & \longrightarrow & 0 \end{array}$$

is fibrant, as fibrations are stable under pullback. Choose a path object

$$\begin{array}{ccc} Y & \xrightarrow{\Delta} & Y \times Y \\ & \searrow t & \nearrow p_0 \times p_1 \\ & & Y' \end{array}$$

with $p_0 \times p_1$ a fibration, t a weak equivalence, and a homotopy

$$\begin{array}{ccc} X & \xrightarrow{K} & Y' \\ & \searrow f_i & \downarrow p_i \\ & & Y \end{array}$$

Taking a fibrant replacement, $Y \xrightarrow{r_Y} RY'$, we obtain a lift

$$\begin{array}{ccc} Y & & \\ \downarrow t & \searrow \Delta & \\ Y' & \xrightarrow{p_0 \times p_1} & Y \times Y \\ \downarrow r_Y & \nearrow \exists q_0 \times q_1 & \downarrow \\ RY' & \longrightarrow & 0 \end{array}$$

because $Y \times Y$ is fibrant and r_Y is a trivial cofibration, making RY' a path object. We also obtain a homotopy $r_Y \circ K$ since the diagrams

$$\begin{array}{ccccc} X & \xrightarrow{K} & Y' & \xrightarrow{r_Y} & RY' \\ & \searrow & \searrow p_0 \times p_1 & \searrow q_0 \times q_1 & \downarrow \\ & & & & Y \times Y \\ & \searrow f_i & \searrow p_i & & \downarrow \pi_i \\ & & & & Y \end{array}$$

commute for $i = 0, 1$. ■

Lemma 5. *Let \mathcal{C} and \mathcal{D} be model categories. Given a Quillen adjunction*

$$\mathcal{C} \begin{array}{c} \xrightarrow{\mathcal{F}} \\ \xleftarrow{\mathcal{G}} \end{array} \mathcal{D}$$

we obtain a derived adjunction

$$\mathrm{Ho}(\mathcal{C}) \begin{array}{c} \xrightarrow{L\mathcal{F}} \\ \xleftarrow{R\mathcal{G}} \end{array} \mathrm{Ho}(\mathcal{D})$$

Proof. Denote by ε and η the unit and counit of adjunction, respectively, and by $[-, -]_{\mathcal{C}}$, $[-, -]_{\mathcal{D}}$ the morphisms of the homotopy category. Also, recall that the isomorphism of adjunction is given by the

morphisms

$$\mathcal{D}(\mathcal{F}X, Y) \xrightarrow{\mathcal{G}(\mathcal{F}X, Y)} \mathcal{C}(\mathcal{G}\mathcal{F}X, \mathcal{G}Y) \xrightarrow{h_{\mathcal{G}Y}(\varepsilon_X)} \mathcal{C}(X, \mathcal{G}Y)$$

and

$$\mathcal{C}(X, \mathcal{G}Y) \xrightarrow{\mathcal{F}(X, \mathcal{G}Y)} \mathcal{D}(\mathcal{F}X, \mathcal{F}\mathcal{G}Y) \xrightarrow{h^{\mathcal{F}X}(\eta_Y)} \mathcal{D}(\mathcal{F}X, Y)$$

We first observe that we have natural isomorphisms

$$[L\mathcal{F}X, Y]_{\mathcal{D}} \cong \mathcal{D}(\mathcal{F}QX, RY)/\sim \text{ and } [X, R\mathcal{G}Y]_{\mathcal{C}} \cong \mathcal{C}(QX, \mathcal{G}RY)/\sim .$$

This reduces the problem to showing that the isomorphism of adjunction both preserves and reflects homotopies between cofibrant objects of \mathcal{C} and fibrant objects of \mathcal{D} , for then we can see that for any object X of \mathcal{C} and any object Y of \mathcal{D} , the isomorphism of adjunction descends to a well-defined isomorphism

$$\begin{array}{ccc} \mathcal{D}(\mathcal{F}QX, RY) & \xrightarrow{\sim} & \mathcal{C}(QX, \mathcal{G}RY) \\ \downarrow & & \downarrow \\ \mathcal{D}(\mathcal{F}QX, RY)/\sim & \longrightarrow & \mathcal{C}(QX, \mathcal{G}RY)/\sim \end{array}$$

Towards that end, assume that X is a cofibrant object of \mathcal{C} and Y is a fibrant object of \mathcal{D} . Given $f_0 \sim f_1 \in \mathcal{D}(\mathcal{F}X, Y)$, choose a homotopy from a fibrant path object.

$$\begin{array}{ccc} Y & \xrightarrow{\Delta} & Y \times Y \\ & \searrow t & \nearrow p_0 \times p_1 \\ & & Y' \end{array}$$

with $p_0 \times p_1$ a fibration, t a weak equivalence, and a homotopy

$$\begin{array}{ccc} \mathcal{F}X & \xrightarrow{K} & Y' \\ & \searrow f_i & \downarrow p_i \\ & & Y \end{array}$$

Since \mathcal{G} is right Quillen it preserves products, fibrant objects, weak equivalences between fibrant objects, and (trivial) fibrations, hence we obtain a path object of \mathcal{C}

$$\begin{array}{ccc} \mathcal{G}Y & \xrightarrow{\mathcal{G}(\Delta)} & \mathcal{G}(Y \times Y) \cong \mathcal{G}(Y) \times \mathcal{G}(Y) \\ & \searrow \mathcal{G}(t) & \nearrow \mathcal{G}(p_0 \times p_1) = \mathcal{G}(p_0) \times \mathcal{G}(p_1) \\ & & \mathcal{G}Y' \end{array}$$

and a homotopy

$$\begin{array}{ccc} X & \xrightarrow{\mathcal{G}(K) \circ \varepsilon_X} & \mathcal{G}Y' \\ & \searrow \mathcal{G}(f_i) \circ \varepsilon_X & \downarrow \mathcal{G}(p_i) \\ & & \mathcal{G}Y \end{array}$$

by the commutativity of the naturality square

$$\begin{array}{ccc} \mathcal{D}(\mathcal{F}X, Y') & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}Y') \\ \downarrow h^{\mathcal{F}X}(p_i) & & \downarrow h^X(\mathcal{G}(p_i)) \\ \mathcal{D}(\mathcal{F}X, Y) & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}Y) \end{array}$$

It's clear that this morphism is surjective and natural, so it remains to show that that it reflects homotopies.

Assume that there is a homotopy $\mathcal{G}(f_0) \circ \varepsilon_X \sim \mathcal{G}(f_1) \circ \varepsilon_X$ in \mathcal{C} . By duality, the argument above implies that we may choose a cofibrant cylinder object

$$\begin{array}{ccc} X \amalg X & \xrightarrow{\nabla} & X \\ & \searrow & \nearrow s \\ & \mathcal{I}_0 \amalg \mathcal{I}_1 & X' \end{array}$$

and a homotopy

$$\begin{array}{ccc} X & \xrightarrow{\iota_i} & X' \\ & \searrow \mathcal{G}(f_i) \circ \varepsilon_X & \downarrow H \\ & & \mathcal{G}Y \end{array}$$

As \mathcal{F} is left Quillen it preserves cofibrant objects, (trivial) cofibrations, and weak equivalences between cofibrant objects by which we obtain a cylinder object in \mathcal{D}

$$\begin{array}{ccc} \mathcal{F}(X \amalg X) \cong \mathcal{F}X \amalg \mathcal{F}X & \xrightarrow{\mathcal{F}(\nabla)} & \mathcal{F}X \\ \mathcal{F}(\mathcal{I}_0 \amalg \mathcal{I}_1) = \mathcal{F}(\mathcal{I}_0) \amalg \mathcal{F}(\mathcal{I}_1) & \searrow & \nearrow \mathcal{F}(s) \\ & & \mathcal{F}X' \end{array}$$

and a homotopy

$$\begin{array}{ccc} \mathcal{F}X & \xrightarrow{\mathcal{F}(\iota_i)} & \mathcal{F}X' \\ & \searrow \eta_Y \circ \mathcal{F}(\mathcal{G}(f_i) \circ \varepsilon_X) & \downarrow \eta_Y \circ \mathcal{F}(H) \\ & & Y \end{array}$$

and we note that

$$\eta_Y \circ \mathcal{F} \circ \mathcal{G}(f_i) \circ \mathcal{F}(\varepsilon_X) = f_i$$

implies that we have a homotopy $f_0 \sim f_1$, as desired. Therefore the induced map is injective, hence an isomorphism. ■

4. QUILLEN EQUIVALENCES

Definition 5. A Quillen adjunction $(\mathcal{F}, \mathcal{G}, \varphi): \mathcal{C} \rightarrow \mathcal{D}$ is a *Quillen equivalence* if and only if, for all cofibrant X in \mathcal{C} and fibrant Y in \mathcal{D} , a map $f: \mathcal{F}X \rightarrow Y$ is a weak equivalence in \mathcal{D} if and only if $\varphi(f): X \rightarrow \mathcal{G}Y$ is a weak equivalence in \mathcal{C} .

Remark 4. Note that a Quillen equivalence is *not* always an equivalence of categories. This could be thought of as a weak equivalence of model categories.

Proposition 5. Let $(\mathcal{F}, \mathcal{G}, \varphi): \mathcal{C} \rightarrow \mathcal{D}$ be a Quillen adjunction with unit $\varepsilon: \text{id}_{\mathcal{C}} \rightarrow \mathcal{G} \circ \mathcal{F}$ and counit $\eta: \mathcal{F} \circ \mathcal{G} \rightarrow \text{id}_{\mathcal{D}}$. The following are equivalent:

- (a) $(\mathcal{F}, \mathcal{G}, \varphi)$ is a Quillen equivalence,
- (b) The composition $X \xrightarrow{\varepsilon_X} \mathcal{G}\mathcal{F}X \xrightarrow{\mathcal{G}(r_{\mathcal{F}X})} \mathcal{G}R\mathcal{F}X$ is a weak equivalence for all cofibrant X , and the composition $\mathcal{F}Q\mathcal{G}Y \xrightarrow{\mathcal{F}(q_{\mathcal{G}Y})} \mathcal{F}QY \xrightarrow{\eta_Y} Y$ is a weak equivalence for all fibrant Y ,
- (c) The derived adjunction is an equivalence of categories.

Proof. (a) \Rightarrow (b)

Assume X is cofibrant in \mathcal{C} and Y is fibrant in \mathcal{D} . Note that $r_{\mathcal{F}X}$ and $q_{\mathcal{G}Y}$ are both weak equivalences. From the diagrams

$$\begin{array}{ccccc}
 & & \varphi & & \\
 & \swarrow & & \searrow & \\
 \mathcal{D}(\mathcal{F}X, R\mathcal{F}X) & \xrightarrow{\mathcal{F}(X, R\mathcal{F}X)} & \mathcal{C}(\mathcal{G}\mathcal{F}X, \mathcal{G}R\mathcal{F}X) & \xrightarrow{h_{\mathcal{G}R\mathcal{F}X}(\varepsilon_X)} & \mathcal{C}(X, \mathcal{G}R\mathcal{F}X) \\
 & & & & \\
 r_{\mathcal{F}X} & \dashv \longrightarrow & \mathcal{G}(r_{\mathcal{F}X}) & \dashv \longrightarrow & \mathcal{G}(r_{\mathcal{F}X}) \circ \varepsilon_X
 \end{array}$$

and

$$\begin{array}{ccccc}
 & & \varphi^{-1} & & \\
 & \swarrow & & \searrow & \\
 \mathcal{C}(Q\mathcal{G}Y, \mathcal{G}Y) & \longrightarrow & \mathcal{D}(\mathcal{F}Q\mathcal{G}Y, \mathcal{F}\mathcal{G}Y) & \longrightarrow & \mathcal{D}(\mathcal{F}\mathcal{G}Y, Y) \\
 & & & & \\
 q_{\mathcal{G}Y} & \dashv \longrightarrow & \mathcal{F}(q_{\mathcal{G}Y}) & \dashv \longrightarrow & \eta_Y \circ \mathcal{F}(q_{\mathcal{G}Y})
 \end{array}$$

we see $\varphi(r_{\mathcal{F}X}) = \mathcal{G}(r_{\mathcal{F}X} \circ \varepsilon_X)$ and $\varphi(\eta_Y \circ \mathcal{F}(q_{\mathcal{G}Y})) = q_{\mathcal{G}Y}$ imply $\mathcal{G}(r_{\mathcal{F}X}) \circ \varepsilon_X$ and $\eta_Y \circ \mathcal{F}(q_{\mathcal{G}Y})$ are weak equivalences by the definition of a Quillen equivalence.

(b) \Rightarrow (c)

It suffices to show that the unit and counit of the adjunction are both isomorphisms. First assume that X is a cofibrant object of \mathcal{C} and note that $\mathcal{F}X$ is cofibrant in X . Chasing the identity through the diagram

$$\begin{array}{ccc}
 \mathcal{D}(\mathcal{F}X, \mathcal{F}X) & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}\mathcal{F}X) \\
 \downarrow h^{\mathcal{F}X}(r_{\mathcal{F}X}) & & \downarrow h^X(\mathcal{G}(r_{\mathcal{F}X})) \\
 \mathcal{D}(\mathcal{F}X, R\mathcal{F}X) & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}R\mathcal{F}X) \\
 \downarrow & & \downarrow \\
 \mathcal{D}(\mathcal{F}X, R\mathcal{F}X) / \sim & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}R\mathcal{F}X) / \sim \\
 \downarrow \wr & & \downarrow \wr \\
 [L\mathcal{F}X, L\mathcal{F}X]_{\mathcal{D}} & \xrightarrow{\sim} & [X, R\mathcal{G} \circ L\mathcal{F}X]_{\mathcal{C}}
 \end{array}$$

we see that the image in the bottom right hand corner is the isomorphism

$$X \xrightarrow{\varepsilon_X} \mathcal{G}\mathcal{F}X \xrightarrow{\mathcal{G}(r_{\mathcal{F}X})} \mathcal{G}R\mathcal{F}X$$

and thus for arbitrary X we obtain the isomorphism

$$X \xrightarrow{q_X^{-1}} QX \xrightarrow{\varepsilon_{QX}} \mathcal{G}\mathcal{F}QX \xrightarrow{\mathcal{G}(r_{\mathcal{F}QX})} \mathcal{G}R\mathcal{F}QX$$

The proof that

$$\mathcal{F}Q\mathcal{G}RY \xrightarrow{\mathcal{F}(q_{\mathcal{G}RY})} \mathcal{F}\mathcal{G}RY \xrightarrow{\eta_{RY}} RY \xrightarrow{r_Y^{-1}} Y$$

is an isomorphism follows similarly. One need only check that the diagrams from part (ii) of Proposition 1 commute.

(c) \Rightarrow (a)

This follows formally from part (iii) of Proposition 4. Let X be a cofibrant object of \mathcal{C} , Y a cofibrant object of \mathcal{D} , and $\varepsilon: \text{id}_{\text{Ho}(\mathcal{C})} \rightarrow R\mathcal{G} \circ L\mathcal{F}$ the unit of the derived adjunction. Note that by assumption both ε is a natural isomorphism.

Given a morphism $f \in \mathcal{D}(\mathcal{F}X, Y)$ denote its image under the isomorphism of adjunction $\tilde{f} \in \mathcal{C}(X, \mathcal{G}Y)$. We have a commutative diagram

$$\begin{array}{ccc} \mathcal{D}(\mathcal{F}X, Y) & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}Y) \\ \downarrow & & \downarrow \\ \mathcal{D}(\mathcal{F}X, Y)/\sim & \xrightarrow{\sim} & \mathcal{C}(X, \mathcal{G}Y)/\sim \\ \downarrow \sim & & \downarrow \sim \\ [L\mathcal{F}X, Y]_{\mathcal{D}} & \xrightarrow{\sim} & [X, R\mathcal{G}Y]_{\mathcal{C}} \end{array}$$

and note that f is a weak equivalence if and only if its image in $\text{Ho}(\mathcal{D})$ is an isomorphism. The image of f in $[X, R\mathcal{G}Y]$ is $R\mathcal{G}(f) \circ \varepsilon_X$, which is an isomorphism if and only if $R\mathcal{G}(f)$ is an isomorphism. Noting that $R\mathcal{G}$ reflects isomorphisms because it is fully faithful, we see that $R\mathcal{G}(f)$ is an isomorphism if and only if the image of f in $\text{Ho}(\mathcal{D})$ is an isomorphism. Since $R\mathcal{G}(f) \circ \varepsilon_X$ is the image of \tilde{f} in $\text{Ho}(\mathcal{C})$, it now follows that f is a weak equivalence if and only if \tilde{f} is a weak equivalence. Therefore the adjoint pair \mathcal{F}, \mathcal{G} form a Quillen equivalence. \blacksquare

5. APPENDIX

Lemma 6. *If X is a cofibrant object of \mathcal{C} , then $q_X: QX \rightarrow X$ is an isomorphism. Dually, if Y is a fibrant object of \mathcal{C} , then $r_Y: Y \rightarrow RY$ is an isomorphism.*

Proof. We have

$$\begin{array}{ccc} 0 & \xrightarrow{\quad} & X \\ & \searrow & \nearrow q_X \\ & & QX \end{array}$$

with q_X a trivial fibration, which gives a lift

$$\begin{array}{ccc}
 0 & \longrightarrow & QX \\
 \downarrow & \nearrow \exists h & \downarrow q_X \\
 X & \xrightarrow{\text{id}_X} & X
 \end{array}$$

because X is cofibrant. We note that h is a weak equivalence by 2-out-of-3. By the functorial factorization we get

$$\begin{array}{ccc}
 X & \xrightarrow{h} & QX \\
 \searrow \alpha(h) & & \nearrow \beta(h) \\
 & X' &
 \end{array}$$

giving a retract

$$\begin{array}{ccccc}
 & & \text{id}_X & & \\
 & & \curvearrowright & & \\
 X & \xrightarrow{\alpha(h)} & X' & \xrightarrow{q_X \circ \beta(h)} & X \\
 \downarrow h & & \downarrow \beta(h) & & \downarrow h \\
 QX & \xrightarrow{\text{id}_{QX}} & QX & \xrightarrow{\text{id}_{QX}} & QX \\
 & & \text{id}_{QX} & & \\
 & & \curvearrowleft & &
 \end{array}$$

and hence h is a trivial fibration. This now gives a lift

$$\begin{array}{ccc}
 0 & \longrightarrow & X \\
 \downarrow & \nearrow \exists h' & \downarrow h \\
 QX & \xrightarrow{\text{id}_{QX}} & QX
 \end{array}$$

and so we see

$$q_X = q_X \circ \text{id}_{QX} = q_X \circ h \circ h' = \text{id}_X \circ h' = h'.$$

■